

- Static Partition
- Dynamic Partition
 - First Fit
 - Best Fit
- Relocatable Dynamic Arrays
- Paging

Relocatable Dynamic Arrays (continued...)

Adv ① Eliminates Fragmentation

② The content of the Job's address space do not have to be altered.

③ A job does not have to be brought back to the same space when it is swapped in/out.

Disadv ① Moving a job requires read and write - consumes time.

② Job's logical address space must be less or equal to the physical address space

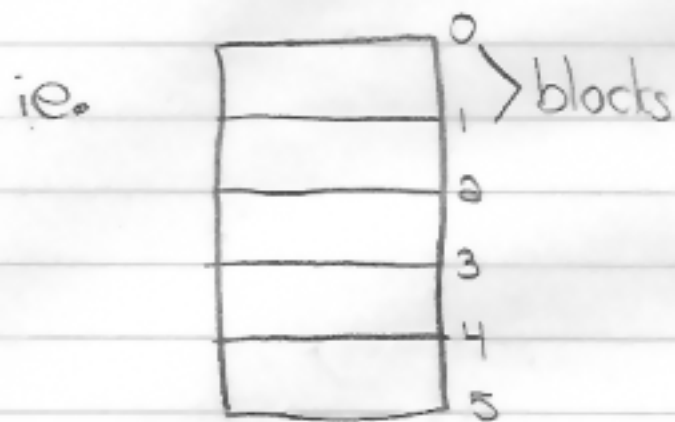
③ "Fossil" parts of a job remains in the memory - hence wasteful

→ Paging memory management schemes resolve the remaining disadvantages.
(Virtual Memory)

Ideas

I The physical address space is (viewed as) a collection of EQUAL SIZE in continuous Blocks

a) A physical address a_p is viewed as a pair (b, l) where
 $b \leftarrow$ block number
 $l \leftarrow$ line number



b) If the block size = 2^m then
 (i) $a_p = b * 2^m + l$

(ii) $b = a_p \text{ DIV } 2^m$
 $l = a_p \text{ MOD } 2^m$

eg. block size = 2^5
 $a_p = \boxed{100100} \boxed{01011}$
 $\quad \quad \quad b \quad \quad \quad l$

II The logical address space is also viewed as a collection of equal sized and continuous regions called pages.

III Page size = Block size

a) A logical address a_e is expressed as a pair (p, L) where
 $P \propto$ page number
 $L \propto$ Line number

b) (i) $a_e = p * 2^m + L$
 (ii) $p = a_e \text{ DIV } 2^m$
 $L = a_e \text{ MOD } 2^m$

IV ANY PAGE can be mapped on any block

V All the pages do not have to be in the physical memory at the same time. Only active pages can be in the memory at a time.

VI A job's logical address space is continuous but its corresponding physical address may not be continuous.

Page Map Table (PMT)

page	block	in memory flag
0	1	Y
1	3	Y
2	17	Y
...
n	...	N

PMT of Job i

VII The mapping between the logical address (A_e) (binding) and its corresponding physical address (A_p) is as follows:

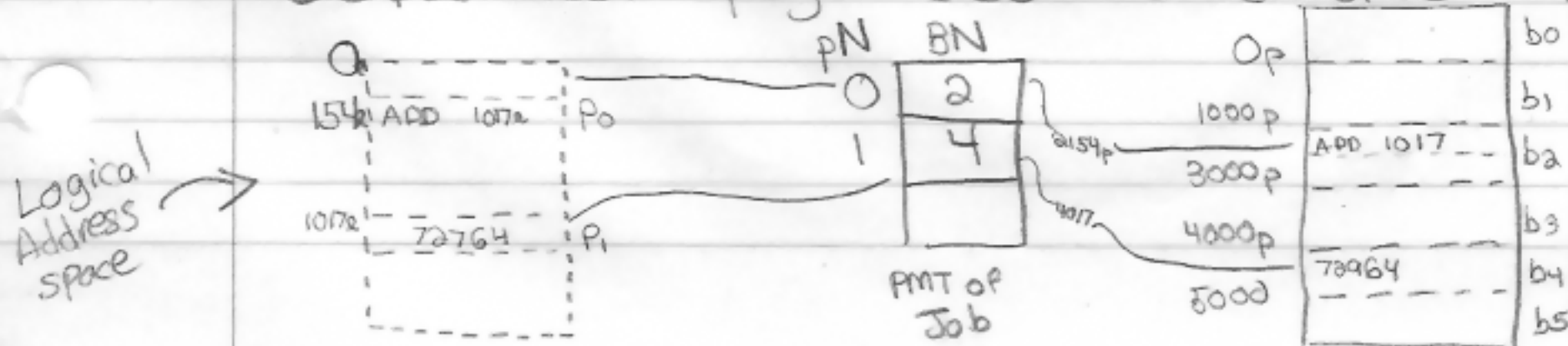
$$A_e = (p, L) \text{ where } p = A_e \text{ DIV } 2^m$$

$$L = A_e \text{ MOD } 2^m$$

$$\therefore b = \text{PMT}(p)$$

$$\therefore A_p = b + 2^m + L$$

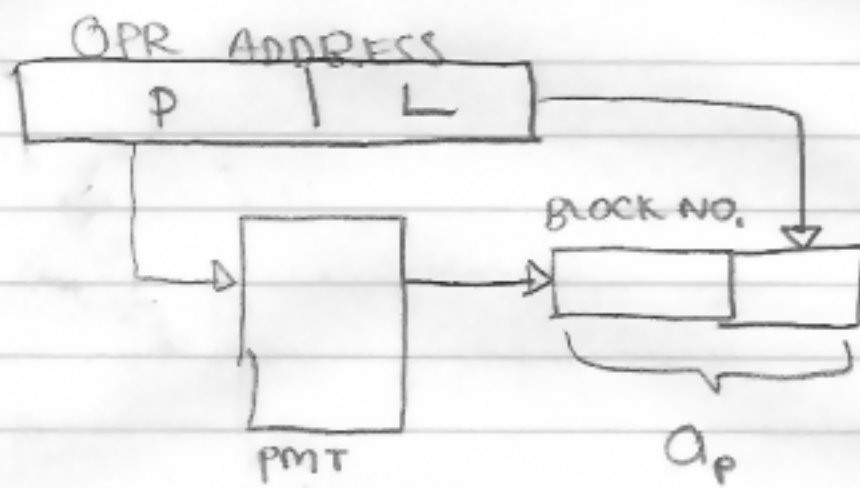
Example Let page size = block size = 1K



$$\text{OPR} = 1017$$

$$p = 1 \Rightarrow \text{PMT}(1) = 4 = b \quad \left. \begin{array}{l} \\ \Rightarrow L = 17 \end{array} \right\} a_p = 4017_p$$

Hardware Support



Exam comprehensive

→ memory management question (1)