

# CS 2742 (Logic in Computer Science) – Fall 2008

## Lecture 4

Antonina Kolokolova

September 18, 2009

Absorption laws:  $p \vee (p \wedge q) \iff p \wedge (p \vee q) \iff p$ .

### 1.1 Conditional statements

Conditional statements are ones of the form “if  $p$  then  $q$ ”,  $p \rightarrow q$ . Recall that we logically define  $p \rightarrow q \iff (\neg p \vee q)$ .

We use the following terminology when talking about conditional statements:

- 1) *Contrapositive* of  $p \rightarrow q$  is  $\neg q \rightarrow \neg p$ . True whenever the original implication is.

*Proof.* Recall that  $(p \rightarrow q) \iff (\neg p \vee q)$ . Now,

$$\begin{aligned} & \neg q \rightarrow \neg p \\ \iff & (\neg \neg q \vee \neg p) && \text{Definition of } \rightarrow \\ \iff & (q \vee \neg p) && \text{Double negation} \\ \iff & (\neg p \vee q) && \text{Commutativity} \iff (p \rightarrow q) && \text{Definition of } \rightarrow \end{aligned}$$

Thus, a contrapositive of an if-then statement is logically equivalent to the original statement.  $\square$

- 2) *Converse* and *inverse*:  $q \rightarrow p$  and  $\neg p \rightarrow \neg q$ . Contrapositives of each other, can have a different truth value from  $p \rightarrow q$ .
- 3) *Sufficient* condition:  $p$  is sufficient for  $q$  if  $p \rightarrow q$ . The “if” part of if-then.
- 4) *Necessary* condition:  $p$  is necessary for  $q$  if  $\neg p \rightarrow \neg q$ , that is,  $q \rightarrow p$ . The “then” part of “if-then”.

5) If and only if ( $p$  iff  $q$ ,  $p \leftrightarrow q$ ) means  $(p \rightarrow q) \wedge (q \rightarrow p)$ .

**Example 1.** Consider the sentence “if  $n$  is divisible by 4, then  $n$  is divisible by 2” (we will use the notation  $n|4$  to mean  $n$  is divisible by 4). This is an if-then statement. Its contrapositive is “if  $n \nmid 2$  then  $n \nmid 4$ . That is, if  $n$  is an odd number then it is definitely not divisible by 4. So  $n|4$  is sufficient for  $n|2$  (if  $n$  is divisible by 4, it is sufficient for  $n$  to be divisible by 2). On the other hand,  $n|2$  is necessary for  $n|4$ .

- 1) *Direct proof:* show that if  $p$  is true directly.
- 2) *Proof by contrapositive:* instead of  $p \rightarrow q$  prove  $\neg q \rightarrow \neg p$ .

**Lemma 1.** *If  $n^2$  is even, then  $n$  is even.*

*Proof.* We will show this by showing that if  $n$  is odd, then  $n^2$  is odd. If  $n$  is odd, then  $n = 2k + 1$  for some  $k$ . Then  $(2k + 1)^2 = 2(2k^2 + 2k) + 1$ , which is an odd number. This proves that if  $n$  is odd, then  $n^2$  is odd, thus proving the contrapositive of if  $n^2$  is even then  $n$  is even, and so proving the statement “if  $n^2$  is even then  $n$  is even” itself.  $\square$

- 3) *Proof by contradiction:* to show that  $p$  is true, show that  $\neg p \rightarrow F$ . It is easy to show that  $(\neg p \rightarrow F)$  is logically equivalent to  $p$ : just note that  $(\neg\neg p \vee F) \iff (p \vee F) \iff p$  by applying the definition of implication followed by the double negation law followed by the identity law.

**Theorem 1.**  *$\sqrt{2}$  is irrational.*

*Proof.* Recall that a number is called *rational* if it can be represented as an (irreducible) fraction of two integers. Assume, for the sake of contradiction, that  $\sqrt{2}$  is rational: that is, there are integers  $m$  and  $n$  which do not have any common divisors  $> 1$  such that  $\sqrt{2} = m/n$ .

- If  $\sqrt{2} = m/n$  then  $(\sqrt{2})^2 = m^2/n^2$ .
- From here,  $2n^2 = m^2$ , which means that  $m^2$  is even.
- By the lemma above, then  $m$  is even, so  $m = 2k$  for some  $k$ .
- Then  $m^2 = 4k^2$ . So  $2n^2 = 4k^2$ , and, dividing by 2,  $n^2 = 2k^2$ . So  $n^2$  is even.
- Using the lemma again, conclude that  $n$  is even.
- So both  $n$  and  $m$  are even, but we assumed that  $m$  and  $n$  do not have a non-trivial common divisor. This is a contradiction.

$\square$

4) *Proof by cases:* to show that  $p$  is true, prove  $(q \rightarrow p) \wedge (\neg q \rightarrow p)$ .

**Lemma 2.** For any natural number  $n$ ,  $\lfloor (n+1)/2 \rfloor = \lceil n/2 \rceil$

Here,  $\lfloor k \rfloor$  is a *floor* of a (real) number, defined to be the largest integer smaller than or equal to  $k$ . For example,  $\lfloor 5/2 \rfloor = 2$ , and  $\lfloor 4/2 \rfloor = 2$ . Similarly, a *ceiling* of a number is the smallest integer larger than or equal to that number. So  $\lceil 5/2 \rceil = 3$  and  $\lceil 4/2 \rceil = 2$ . For an integer, both its floor and its ceiling are equal to that integer itself; for a non-integer, the floor is its rounded-down value and the ceiling is rounded up.

*Proof.* **Case 1:  $n$  is even.** Then  $\lfloor (n+1)/2 \rfloor = n/2 = \lceil n/2 \rceil$ .

**Case 2:  $n$  is odd.** Then  $\lfloor (n+1)/2 \rfloor = (n+1)/2 = \lceil n/2 \rceil$ . □

**Puzzle 4.** A from the island of knights and knaves said: “If I am a knight, then I’ll eat my hat!”. Prove that A will eat his hat.