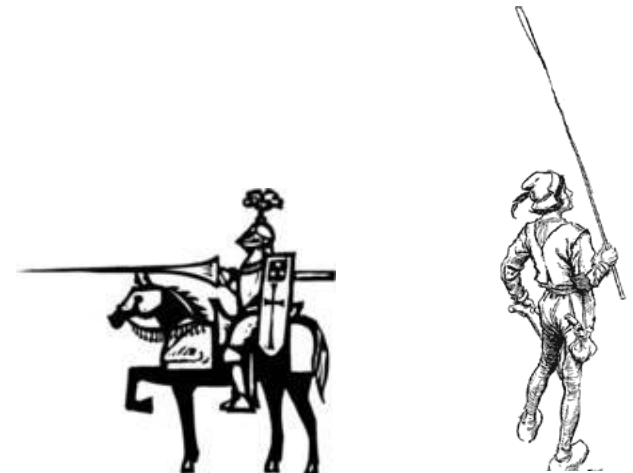


COMP 1002

Intro to Logic for Computer Scientists

Lecture 3



Admin stuff

- Make-up lecture next Monday, Jan 14
- 10am to 12pm in C 3033





Knights and knaves



- On a mystical island, there are two kinds of people: knights and knaves. Knights always tell the truth. Knaves always lie.
- Puzzle 1: You meet two people on the island, Arnold and Bob. Arnold says “Either I am a knave, or Bob is a knight”. Is Arnold a knight or a knave? What about Bob?



Knights and knaves



- Puzzle 1: You meet two people on the island, Arnold and Bob. Arnold says “Either I am a knave, or Bob is a knight”. Is Arnold a knight or a knave? What about Bob?
 - A: Arnold is a knight
 - B: Bob is a knight
 - Formula: $\neg A \vee B$: “Either Arnold is a knave, or Bob is a knight”
 - Want: scenarios where either both A is a knight and the formula is true, or A is a knave and the formula is false. Use “if and only if” notation: $(\neg A \vee B) \leftrightarrow A$. True if both formulas have same value

A	B	$\neg A$	$\neg A \vee B$	$(\neg A \vee B) \leftrightarrow A$
<i>True</i>	<i>True</i>	False	True	True
<i>True</i>	<i>False</i>	False	False	False
<i>False</i>	<i>True</i>	True	True	False
<i>False</i>	<i>False</i>	True	True	False

Special types of sentences

- A sentence that has a satisfying assignment is **satisfiable**.
 - Some row in the truth table ends with *True*.
 - Example: $B \rightarrow A$

A	B	$B \rightarrow A$
<i>True</i>	<i>True</i>	<i>True</i>
<i>True</i>	<i>False</i>	<i>True</i>
<i>False</i>	<i>True</i>	<i>False</i>
<i>False</i>	<i>False</i>	<i>True</i>

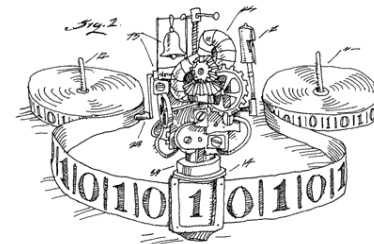
- Sentence is a **contradiction**:
 - All assignments are falsifying.
 - All rows end with *False*.
 - Example: $A \wedge \neg A$

A	$A \wedge \neg A$
<i>True</i>	<i>False</i>
<i>False</i>	<i>False</i>

- Sentence is a **tautology**:
 - All assignments are satisfying
 - All rows end with *True*.
 - Example: $B \rightarrow A \vee B$

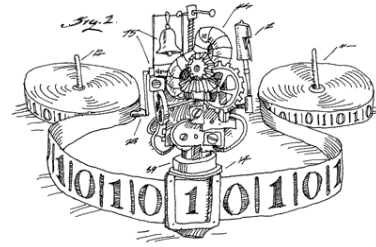
A	B	$A \vee B$	$B \rightarrow A \vee B$
<i>True</i>	<i>True</i>	<i>True</i>	<i>True</i>
<i>True</i>	<i>False</i>	<i>True</i>	<i>True</i>
<i>False</i>	<i>True</i>	<i>True</i>	<i>True</i>
<i>False</i>	<i>False</i>	<i>False</i>	<i>True</i>

Determining formula type



- How long does it take to check if a formula is satisfiable?
 - If somebody gives you a satisfying assignment, then in time roughly the size of the formula.
 - On a m -symbol formula, take time $O(m) = \text{constant} * m$, for some constant depending on the computer/software.
 - What if you don't know a satisfying assignment? How hard it is to find it?
 - Using a truth table: in time $O(m * 2^n)$ on a length m n -variable formula.
 - Is it efficient?...

Complexity of computation

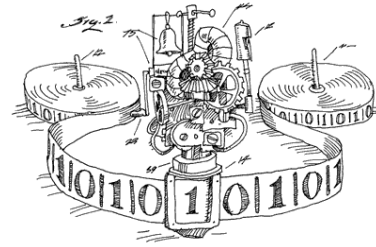


- Would you still consider a problem really solvable if it takes very long time?
 - Say 10^n steps on an n -symbol string?
 - At a billion (10^9) steps per second (~ 1 GHz)?
 - To process a string of length 100...
 - will take $10^{100}/10^9$ seconds, or $\sim 3 \times 10^{72}$ centuries.



- Age of the universe: about 1.38×10^{10} years.
- Atoms in the observable universe: 10^{78} - 10^{82} .

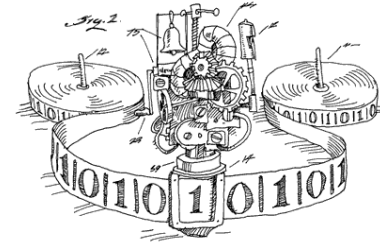
Complexity of computation



- What strings do we work with in real life?
 - A DNA string has 3.2×10^9 base pairs
 - A secure key in crypto: 128-256 bits
 - Number of Walmart transactions per day: 10^6 .
 - URLs searched by Google in 2012: 3×10^{12} .



Determining formula type



- How long does it take to check if a formula is satisfiable?
 - Using a truth table: in time $O(m * 2^n)$ on a length m n -variable formula.
 - Is it efficient?
 - Not really!
 - Formula with 100 variables is already too big!
 - In software verification: millions of variables!
 - Can we do better?



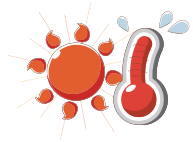
A million-dollar question!



Double negation



- Negation cancels negation
 - $\neg\neg A \equiv A$
 - “I do not disagree with you” = “I agree with you”
- For a human brain, harder to parse a sentence with multiple negations:
 - Alice says: “I refuse to vote against repealing the ban on smoking in public.”
 - Does Alice like smoking in public or hate it?

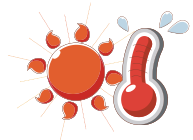




De Morgan's Laws



- Simplifying negated formulas
 - For AND: $\neg (A \wedge B)$ is equivalent to $(\neg A \vee \neg B)$
 - For OR: $\neg (A \vee B) \equiv (\neg A \wedge \neg B)$
- Example:
 - $\neg (\neg A \vee B)$ is $\neg \neg A \wedge \neg B$, same as $A \wedge \neg B$
 - So, since $(A \rightarrow B)$ is equivalent to $(\neg A \vee B)$,
 $\neg(A \rightarrow B)$ is equivalent to $A \wedge \neg B$





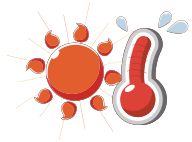
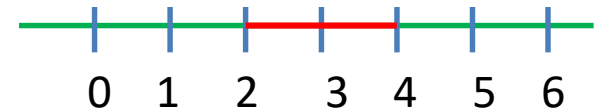
De Morgan's laws: examples



- Let A be “it’s sunny” and B “it’s cold”.
 - “It’s sunny and cold today”! -- No, it’s not!
 - That could mean
 - No, it’s not sunny.
 - No, it’s not cold.
 - No, it’s neither sunny nor cold.
 - In all of these scenarios, “It’s either not sunny or not cold” is true.



- Let A be “ $x < 2$ ”, B be “ $x > 4$ ”.
 - “Either $x < 2$ or $x > 4$ ” – No, it is not!
 - Then $2 \leq x \leq 4$



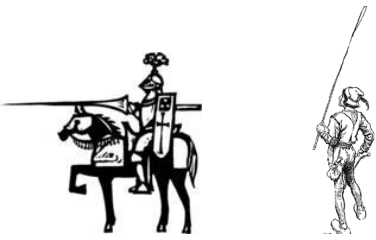
More examples



- Let A be “I play” and B “I win”.
 - $A \rightarrow B$: “If I play, then I win”
 - Equivalent to $\neg A \vee B$: “Either I do not play, or I win”.
- Negation: $\neg(A \rightarrow B)$: “It is not so that if I play then I win”.
 - By de Morgan’s law: $\neg(\neg A \vee B) \equiv (\neg\neg A \wedge \neg B)$
 - By double negation: $(\neg\neg A \wedge \neg B) \equiv (A \wedge \neg B)$
 - So negation of “If I play then I win” is “I play **and** I **don’t** win”.

Longer example of negation

- Start with the outermost connective and keep applying de Morgan's laws and double negation.
- Stop when all negations are on variables.
- $\neg ((A \vee \neg B) \rightarrow (\neg A \wedge C))$
 - $(A \vee \neg B) \wedge \neg(\neg A \wedge C)$ (negating \rightarrow)
 - $(A \vee \neg B) \wedge (\neg\neg A \vee \neg C)$ (de Morgan)
 - $(A \vee \neg B) \wedge (A \vee \neg C)$ (removing $\neg\neg$)



Knights and knaves



- On a mystical island, there are two kinds of people: knights and knaves. Knights always tell the truth. Knaves always lie.
- Puzzle 2: You see three islanders talking to each other, Arnold, Bob and Charlie.
 - You ask Arnold “Are you a knight?”, but can’t hear what he answered.
 - Bob pitches in: “Arnold said that he is a knave!”
 - and Charlie interjects “Don’t believe Bob, he’s lying”.
 - Out of Bob and Charlie, who is a knight/knave?