

COMP 1002

Intro to Logic for Computer Scientists

Lecture 14





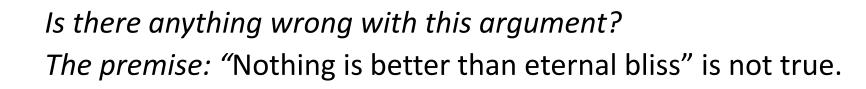


Puzzle: better than nothing

- Nothing is better than eternal bliss
- A burger is better than nothing

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• Therefore, a burger is better than eternal bliss.



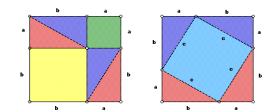






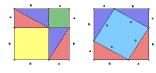


Types of proofs



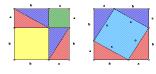
- Direct proof of $\forall x F(x)$
 - Show that F(x) holds for arbitrary x, then use universal generalization.
 - Often, F(x) is of the form $G(x) \rightarrow H(x)$
 - Example: A sum of two even numbers is even.
 - Example: Difference of numbers congruent mod d.
- Proof by cases
 - If can write $\forall x F(x)$ as $\forall x(G_1(x) \lor G_2(x) \lor \cdots \lor G_k(x)) \to H(x)$, prove $(G_1(x) \to H(x)) \land (G_2(x) \to H(x)) \land \cdots \land (G_k(x) \to H(x))$
 - Example: triangle inequality $(|x + y| \le |x| + |y|)$
- Proof by contraposition
 - To prove $\forall x \ G(x) \rightarrow H(x)$, prove $\forall x \neg H(x) \rightarrow \neg G(x)$
 - Example: If square of an integer is even, then this integer is even.
- Proof by contradiction
 - To prove $\forall x F(x)$, prove $\forall x \neg F(x) \rightarrow FALSE$
 - Example: $\sqrt{2}$ is not a rational number.
 - Example: There are infinitely many primes.





Direct proof

- **Direct proof** of $\forall x \in S \ F(x)$: show directly that F(x) holds for arbitrary $n \in S$, then use universal generalization.
 - Universal instantiation: "let n be an arbitrary element of the domain S of $\forall x$ "
 - Show F(n) from axioms, definitions, previous theorems...
 - When F(x) is of the form $G(x) \rightarrow H(x)$, then assume G(n) is true, and from that (and axioms, etc) derive H(n)
 - That proves $G(n) \rightarrow H(n)$
 - Now use universal generalization to conclude that $\forall x F(x)$ is true.



Direct proof

• *Definition* (of even integers):

- An integer n is **even** iff $\exists k \in \mathbb{Z}, n = 2 \cdot k$.

• *Theorem*: Sum of two even integers is even.

 $- \forall x, y \in \mathbb{Z} \ Even(x) \land Even(y) \rightarrow Even(x + y).$

- Proof:
 - Suppose m and n are arbitrary even integers.
 - Universal instantiation.
 - Then $\exists k \in \mathbb{Z}, n = 2k$ and $\exists l \in \mathbb{Z}, m = 2l$.
 - By definition: note different variables.
 - -m + n = 2k + 2l = 2(k + l)
 - By substitution and axioms of theory of integers (algebra).
 - -m + n = 2(k + l), so m + n is even
 - By definition (other direction of iff).
 - Since m and n were arbitrary, therefore, we have shown what we needed: $\forall x, y \in \mathbb{Z}$ $Even(x) \land Even(y) \rightarrow Even(x + y)$.
 - By universal generalization.

Modular arithmetic



- Quotient-remainder theorem: for any integer n and a positive integer d, there exist unique integers q (quotient) and r (reminder) such that: n = dq + r and 0 ≤ r < d
 16 = 3*5+1, 11 = 2*4+3...
- $n \equiv m \pmod{d}$, pronounced "*n* is congruent to *m* mod *d*", means that n and m have the same remainder when divided by d. That is, $n = dq_1 + r$ and $m = dq_2 + r$, for the same r.
 - In some programming languages, there is an operator mod, so you might see "n mod d", which would return r.
 - In Python, it is n % d.
 - $n \equiv m \pmod{d}$ and $m = n \mod{d}$ are not the same:
 - $10 \equiv 16 \pmod{3}$, but $10 \mod 3 = 1$
 - Operator div, "n div d" is sometimes used to compute q.
 - In Python, integer division (or //) does it.

Calendars vs mod





Wednesdays are
 day = 3 (mod 7)

 Wednesdays are day = 4 (mod 7)

Calendars vs mod





Wednesdays are
 day = 3 (mod 7)

 Orange stickers are number = 3 (mod 8)

Modular arithmetic in CS

- Example: day of the week.
 - Feb 1st and Feb 15th are both on Wednesday: $1 \equiv 15 \pmod{7}$
- Hash functions: distribute random data evenly among d memory locations
 - Often take h(k) = k mod p for some prime p. If $k \equiv \ell \pmod{p}$, get a collision.
- Cryptography:
 - Parity checks in credit cards, codes, ISBNs, etc.
 - E.g., look at combination of digits mod 10 to check if a credit card number is valid.
 - Public key crypto, RSA....



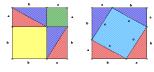
Direct proof example



• Theorem: for all integers n,m and d, where d > 0, if $n \equiv m \pmod{d}$ then there exists an integer k such that n = m + kd

 $- \ \forall x, y, z \ (z > 0 \land x \equiv y \ (mod \ z)) \rightarrow \exists u \ x = y + uz$

- Proof:
 - Let n, m, d be arbitrary integers such that d > 0 and $n \equiv m \pmod{d}$
 - Universal instantiation and assuming the premise
 - Then there are integers q_1, q_2, r with $0 \le r < d$ such that $n = dq_1 + r$ and $m = dq_2 + r$.
 - By the quotient-remainder theorem and definition of congruence.
 - Now, $n-m = (dq_1 + r) (dq_2 + r) = d(q_1 q_2)$
 - Substitution and algebra.
 - Set $k = q_1 q_2$. For this k, n = m + kd. Therefore, $\exists u \ n = m + ud$
 - By existential generalization
 - Since n, m, d were arbitrary integers with d > 0 and $n \equiv m \pmod{d}$, $\forall x, y, z \ (z > 0 \land x \equiv y \pmod{z}) \rightarrow \exists u \ x = y + uz$
 - By universal generalization.



Proof by contraposition

- To prove $\forall x \ G(x) \rightarrow H(x)$, prove its contrapositive $\forall x \neg H(x) \rightarrow \neg G(x)$
 - Universal instantiation: "let n be an arbitrary element of the domain S of ∀x "
 - Suppose that $\neg H(n)$ is true.
 - Derive that $\neg G(n)$ is true.
 - Conclude that $\neg H(n) \rightarrow \neg G(n)$ is true.
 - Now use universal generalization to conclude that
 ∀x F(x) is true.



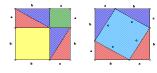
Pigeonhole Principle

- Suppose that nobody in our class carries more than 10 pens.
- There are 108 students in our class.
- Prove that there are at least 2 students in our class who carry the same number of pens.
 - In fact, there are at least 10 who do.
- The Pigeonhole Principle:
 - If there are n pigeons
 - And n-1 pigeonholes
 - Then if every pigeon is in a pigeonhole
 - At least two pigeons sit in the same hole









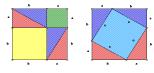
 \Box (Done).

Proof by contraposition.

• Theorem (PigeonHolePrinciple): For any n, if there are n+1 pigeons and n holes, then if every pigeon sits in some hole, then there is a hole with at least two pigeons.

$$\begin{array}{l} - \forall x \in \mathbb{N} \ \left(\forall \ y \in \{1, \dots, x+1\} \exists \ z \in \{1, \dots, x\} \ Sits(y, z) \right) \rightarrow \\ \left(\exists \ u \in \{1, \dots, x+1\} \exists \ v \in \{1, \dots, x+1\} \ \exists w \in \{1, \dots, x\} \\ \left(u \neq v \land Sits(u, w) \land Sits(v, w) \right) \right) \end{array}$$

- Proof:
 - Suppose n is an arbitrary integer.
 - We show the contrapositive: if every hole has at most one pigeon, then some pigeon is not sitting in any hole.
 - If every hole has at most one pigeon, then there are at $\leq 1^*n=n$ pigeons sitting in holes.
 - Then there is (n + 1) n = 1 pigeon that is not sitting in a hole, proving the contrapositive.
 - Therefore, if every pigeon sits in a hole, and there are more than n pigeons, then two pigeons sit in the same hole.
 - By universal generalization, done.



Proof by contraposition.

- *Theorem*: If a square of an integer is even, that integer is even.
 - $\forall x \in \mathbb{Z} Even(x^2) \rightarrow Even(x).$
- Proof:
 - We will show a contrapositive: $\forall x \in \mathbb{Z} \neg Even(x) \rightarrow \neg Even(x^2)$. That is, square of an odd integer is odd.
 - Let n be an arbitrary odd integer. By definition, n = 2k + 1 for some integer k.
 - Then $n^2 = (2k + 1)^2 = 4k^2 + 4k + 1 = 2(2k^2 + 2k) + 1$,
 - So $n^2 = 2m + 1$ for m= $2k^2 + 2k$, thus n^2 is odd by definition.
 - By universal generalization, get $\forall x \in \mathbb{Z} \neg Even(x) \rightarrow \neg Even(x^2)$. Since it is a contrapositive of the original statement, done.

Puzzle: Caesar cipher



- The Roman dictator Julius Caesar encrypted his personal correspondence using the following code.
 - Number letters of the alphabet: A=0, B=1,... Z=25.
 - To encode a message, replace every letter by a letter three positions before that (wrapping).
 - A letter numbered x by a letter numbered x-3 mod 26.
 - For example, F would be replaced by C, and A by X
- Suppose he sent the following message.
 QOBXPROB FK QEB ZXSB
- What does it say?

