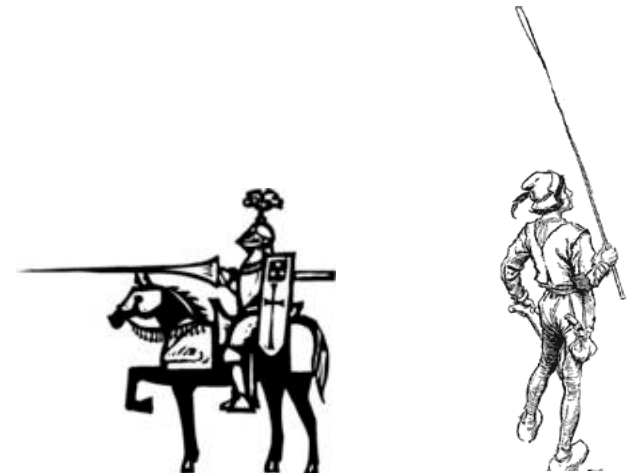


COMP 1002

Intro to Logic for Computer Scientists

Lecture 9



Puzzle 8

- Suppose that nobody in our class carries more than 10 pens.
- There are 70 students in our class.
- Prove that there are at least 2 students in our class who carry the same number of pens.
 - In fact, there are at least 7 who do.



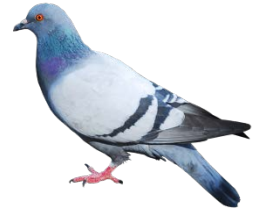
Pigeonhole Principle



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- The Pigeonhole Principle:
 - If there are n pigeons
 - And $n-1$ pigeonholes
 - Then if every pigeon is in a pigeonhole
 - At least two pigeons sit in the same hole



Pigeonhole Principle



- Suppose that nobody in our class carries more than 10 pens. There are 70 students in our class. Prove that there are at least 2 students in our class who carry the same number of pens.
 - In fact, there are at least 7 who do.
- The Pigeonhole Principle:
 - If there are n pigeons and $n-1$ pigeonholes
 - Then if every pigeon is in a pigeonhole
 - At least two pigeons sit in the same hole
- Applying to our problem:
 - $n-1 = 11$ possible numbers of pens (from 0 to 10)
 - Even with $n=12$ people, there would be 2 who have the same number.
 - If there were less than 7, say 6 for each scenario, total would be 66.
 - Note that it does not tell us which number or who these people are!



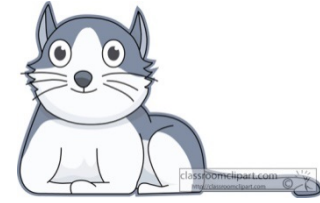
Resolution and Pigeons



- It is not that hard to write the Pigeonhole Principle as a tautology
- But we can prove that resolution has trouble with this kind of reasoning
 - the smallest resolution proof of this tautology is exponential size!
- By contrast, natural deduction (and you!) can figure it out fairly quickly
 - though it is not straightforward.
- The problem is that resolution cannot count.
 - But ability to count makes things harder...



Meow-stery



- One evening there was a cat fight in a family consisting of a mother cat, a father cat, and their son and daughter kittens.
- One of these four cats attacked and bit another!
- One of the cats watched the fight.
- The other one hissed at the fighters.
- These are the things we know for sure:
 - 1. The watcher and the hisser were not of the same sex.
 - 2. The oldest cat and the watcher were not of the same sex.
 - 3. The youngest cat and the victim were not of the same sex.
 - 4. The hissing cat was older than the victim.
 - 5. The father was the oldest of the four.
 - 6. The attacker was not the youngest of the four.
- Which nasty cat was the attacker?



Boolean functions and circuits

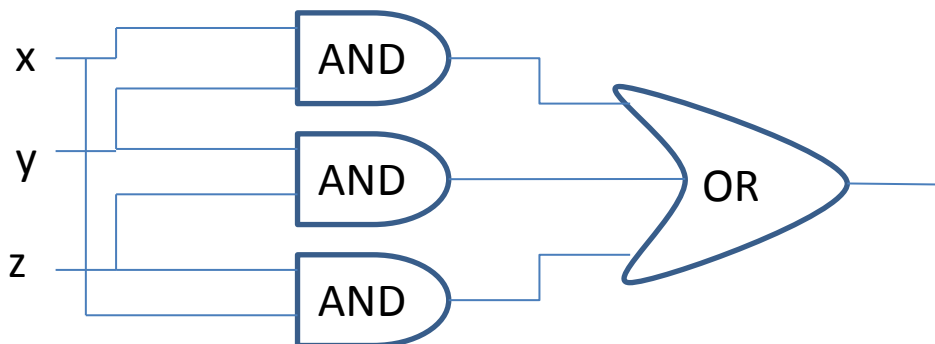


- What is the relation between propositional logic we studied and logic circuits?
 - View a formula as computing a function (called a Boolean function),
 - inputs are values of variables,
 - output is either *true (1)* or *false (0)*.
 - For example, $Majority(x, y, z) = true$ when at least two out of x, y, z are true, and false otherwise.
 - Such a function is fully described by a truth table of its formula (or its circuit: circuits have truth tables too).

Boolean functions and circuits



- What is the relation between propositional logic and logic circuits?
 - So both formulas and circuits “compute” Boolean functions – that is, truth tables.
 - In a circuit, can “reuse” a piece in several places, so a circuit can be smaller than a formula.
 - Still, most circuits are big!
 - *Majority*(x, y, z) is $(x \wedge y) \vee (x \wedge z) \vee (y \wedge z)$



Canonical DNF



- So for every formula, there is a unique canonical CNF (and a truth table, and a Boolean function).
- And for every possible truth table (a Boolean function), there is a formula (the canonical CNF).
- A negation of a CNF is an OR of ANDs of literals. It is called a **DNF (disjunctive normal form)**.
 - To make a canonical DNF from a truth table:
 - take all satisfying assignments.
 - Write each as an AND of literals, as before.
 - Then take an OR of these ANDs.

Complete set of connectives



- CNFs only have \neg, \vee, \wedge , yet any formula can be converted into a CNF
 - Any truth table can be coded as a CNF
- Call a set of connectives which can be used to express any formula **a complete set of connectives**.
 - In fact, \neg, \vee is already complete. So is \neg, \wedge .
 - By DeMorgan, $(A \vee B) \equiv \neg(\neg A \wedge \neg B)$ No need for \vee !
 - But \wedge, \vee is not: cannot do \neg with just \wedge, \vee .
 - Because when both inputs have the same value, both \wedge, \vee leave them unchanged.

Complete set of connectives



- How many connectives is enough?
 - Just one: NAND (NotAND), also called the Sheffer stroke, written as $|$

$$\neg A \equiv A | A$$

$$\begin{aligned} A \vee B &\equiv \neg(\neg A \wedge \neg B) \\ &\equiv (\neg A | \neg B) \\ &\equiv ((A | A) | (B | B)) \end{aligned}$$

A	B	A B
True	True	False
True	False	True
False	True	True
False	False	True

- In practice, most often stick to \wedge, \vee, \neg

Puzzle 9



- Susan is 28 years old, single, outspoken, and very bright. She majored in philosophy. As a student she was deeply concerned with issues of discrimination and social justice and also participated in anti-nuke demonstrations.

Please rank the following possibilities by how likely they are. List them from least likely to most likely. Susan is:

1. a kindergarden teacher
2. works in a bookstore and takes yoga classes
3. an active feminist
4. a psychiatric social worker
5. a member of an outdoors club
6. a bank teller
7. an insurance salesperson
8. a bank teller and an active feminist

