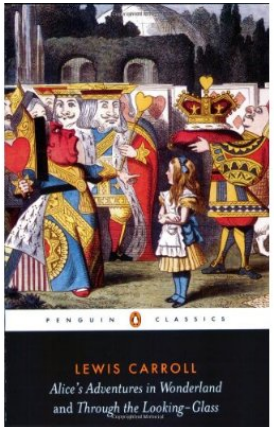
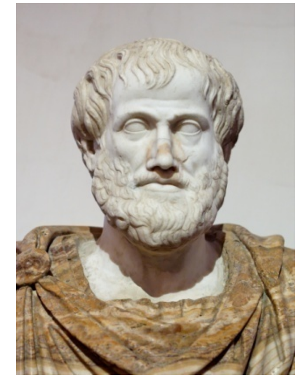


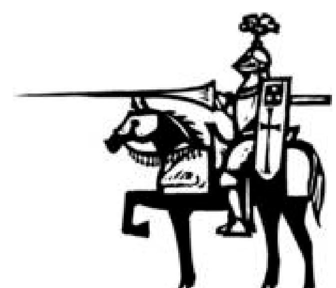


COMP 1002

Logic for Computer Scientists

Lecture 19

B 5 2 J



Admin stuff

- **Midterm March 2nd.**
  - Closed-book.
  - Covers lectures 1 to 16.
    - Mainly labs 1, 2, 3, 4 and assignments 1 and 2.
  - Study guide posted to help you study
    - **not** to bring to the midterm itself.
  - Sample midterm is posted.
    - I will not be posting solutions, but will be happy to give you feedback on yours.

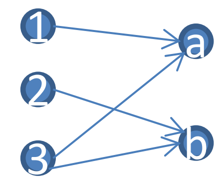
- Assignments 1 and 2 are marked.
    - Except for several assignments 2, will be ready in a couple of hours.
  - Read your feedback on D2L
  - Let me know as soon as possible if you have questions about your mark.

- Lab 5 moved from the week before the break to this Wednesday.

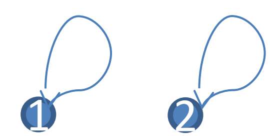
Relations

- **A relation** is a subset of a Cartesian product of sets.
  - If of two sets (set of pairs), call it a **binary** relation.
  - Of 3 sets (set of triples), **ternary**.   Of k sets (set of tuples), **k-ary**

  - A={1,2,3}, B={a,b}
    - $A \times B = \{(1,a),(1,b),(2,a),(2,b),(3,a),(3,b)\}$
    - R = {(1,a), (2,b),(3,a), (3,b)} is a relation. So is R={(1,b)}.
  - A={1,2},
    - $A \times A\ = \{(1,1),(1,2),(2,1),(2,2)\}$
    - R={(1,1), (2,2)} (all pairs (x,y) where x=y)
    - R={(1,1),(1,2),(2,2)} (all pairs (x,y) where $x \leq y$).
  - A=PEOPLE
    - COUPLES ={(x,y) | Loves(x,y)}
    - PARENTS ={(x,y) | Parent(x,y)}
  - A=PEOPLE, B=DOGS, C=PLACES
    - WALKS = {(x,y,z) | x walks y in z}
      - Jane walks Buddy in Bannerman park.

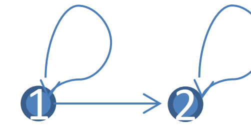


Graph of R (bipartite)



Graph of {(1,1),(2,2)}

Types of binary relations



R_1

- A binary relation $R \subseteq A \times A$ is
  - **Reflexive** if $\forall x \in A, R(x,x)$
    - Every x is related to itself.
    - E.g. A={1,2}, R_1 = { (1,1), (2,2), (1,2)}
    - On A = $\mathbb{Z}$, $R_2 = \{(x,y) | x = y\}$ is reflexive
    - But not $R_3 = \{(x,y) |\ x < y\}$
  - **Symmetric** if $\forall x, y \in A,\ (x,y) \in R \leftrightarrow (y,x) \in R$
    - R_1 and R_3 above are not symmetric. R_2 is.
    - A = $\mathbb{Z}$, $R_4 = \{(x,y) | x \equiv y \ mod\ 3\}$ is symmetric.
  - **Transitive** if $\forall x, y, z \in A,\ (x,y) \in R \wedge (y,z) \in R \rightarrow (x,z) \in R$
    - R_1, R_2, R_3, R_4 are all transitive.
    - $R_5 = \{(x,y) | x, y \in \mathbb{Z} \wedge x + 1 = y\}$ is not transitive.
    - PARENT = $\{(x,y) | x, y \in PEOPLE \wedge x\ \ is\ a\ parent\ of\ y\}$ is not.
    - A **transitive closure** of a relation R is a relation $R^* = \{(x,z) |\ \exists k \in \mathbb{N}\ \ \exists y_0, \dots, y_k \in A\ \ (x = y_0 \wedge z = y_k \wedge \forall i \in \{0, \dots, k-1\}\ R(y_i, y_{i+1})\}$
      - That is, can get from x to z following R arrows.

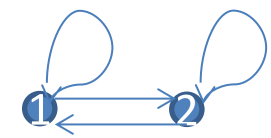
Types of binary relations

Graph of {(1,2)}

- A binary relation $R \subseteq A \times A$ is
  - **Anti-reflexive** if $\forall x \in A, \neg R(x,x)$
    - R can be neither reflexive nor anti-reflexive.
    - E.g. A={1,2}, $R_6 = \{(1,2)\}$
      - but not $R_1 = \{(1,1), (2,2), (1,2)\}$ (reflexive)
      - nor $R_7 = \{(1,1), (1,2)\}$ (neither)
    - For $A = \mathbb{Z}$, not $R_2 = \{(x,y) | x = y\}$
      - Nor $R_4 = \{(x,y) | x \equiv y \; mod \; 3\}$
    - But $R_3 = \{(x,y) | \; x < y\}$ is anti-reflexive.
      - So are $R_5 = \{(x,y) \in \mathbb{Z} \times \mathbb{Z} \; | \; x + 1 = y\}$
      - And PARENT = $\{(x,y) \in PEOPLE \times PEOPLE \; | x \; \; is \; a \; parent \; of \; y\}$
  - **Anti-symmetric** if $\forall x, y \in A, (x,y) \in R \wedge (y,x) \in R \rightarrow x = y$
    - $R_1, R_3, R_5, R_6, R_7, PARENT$ are anti-symmetric. R_4 is not.
    - R_2 is both symmetric and anti-symmetric.
    - $R_8 = \{(1,2), (2,1), (1,3)\}$ is neither symmetric nor anti-symmetric.

Equivalence

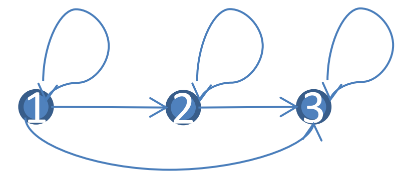
- A binary relation $R \subseteq A \times A$ is an **equivalence** if R is reflexive, symmetric and transitive.
  - E.g. A={1,2}, $R = \{(1,1), (2,2)\}$ or $R = A \times A$
  - Not $R_1 = \{ (1,1), (2,2), (1,2)\}$ nor $R_3 = \{(x, y) | \ x < y\}$
  - On A = $\mathbb{Z}$, $R_2 = \{(x, y) | x = y\}$ is an equivalence
  - So is $R_4 = \{(x, y) | x \equiv y \ mod\ 3 \}$
    - Reflexive: $\forall x \in \mathbb{Z},\ x \equiv x \ mod\ 3$
    - Symmetric: $\forall x, y \in \mathbb{Z},\ x \equiv y \ \mathrm{mod}\ 3 \ \rightarrow y \equiv x \ mod\ 3$
    - Transitive: $\forall x, y, z \in \mathbb{Z},\ x \equiv y \ \mathrm{mod}\ 3 \wedge y \equiv z \ mod\ 3 \rightarrow x \equiv z \ mod\ 3$



- An equivalence relation partitions A into **equivalence classes**:
  - Intersection of any two equivalence classes is ∅
  - Union of all equivalence classes is A.
  - R_4: $\mathbb{Z} = \{x \ | x \equiv 0 \ mod\ 3\} \cup \{\mathrm{x} \ | \ \mathrm{x} \equiv 1 \ mod\ 3\} \cup \{x \ | x \equiv 2 \ mod\ 3\}$
  - $R = A \times A$ gives rise to a single equivalence class. $R = \{(1,1), (2,2)\}$ to two.

Partial and total orders

- A binary relation $R \subseteq A \times A$ is an **order** if R is reflexive, anti-symmetric and transitive.
  - R is a **total order** if $\forall x, y \in A \;\; R(x,y) \vee R(y,x)$
    - That is, every two elements of A are related.
    - E.g. $R_1 = \{(x,y) | x, y \in \mathbb{Z} \wedge x \leq y\}$ is a total order.
    - So is alphabetical order of English words.
    - But not $R_2 = \{(x,y) | x, y \in \mathbb{Z} \wedge x < y\}$
      - not reflexive, so not an order.



  - Otherwise, R is a **partial order**.
    - $SUBSETS = \{(A,B) \mid A, B \text{ are sets} \wedge \; A \subseteq B\}$ is a partial order.
      - Reflexive: $\forall A, \;\; A \subseteq A$
      - Anti-symmetric: $\forall A, B \;\; A \subseteq B \wedge B \subseteq A \rightarrow A = B$
      - Transitive: $\forall A, B, C \;\; A \subseteq B \wedge B \subseteq C \rightarrow A \subseteq C$
      - Not total: if A ={1,2} and B ={1,3}, then neither $A \subseteq B$ nor $B \subseteq A$

    - $DIVISORS = \{(x,y) | \; x, y \in \mathbb{N} \wedge x, y \geq 2 \; \wedge \exists z \in \mathbb{N} \;\; y = z \cdot x\}$ is a partial order.
    - PARENT is not an order. But ANCESTOR would be, if defined so that each person is an ancestor of themselves. It is a partial order.
- An order may have **minimal** and **maximal** elements (maybe multiple)
  - $x \in A$ is minimal in R if $\forall y \in A \;\; y \neq x \rightarrow \; \neg R(y,x)$
    - and maximal if $\forall y \in A \; y \neq x \rightarrow \neg R(x,y)$
  - ∅ is minimal in SUBSETS (its unique minimum); universe is maximal (its unique maximum).
  - All primes are minimal in DIVISORS, and there are no maximal elements.

Puzzle: coins

- A not-too-far-away country recently got rid of a penny coin, and now everything needs to be rounded to the nearest multiple of 5 cents…
  - Suppose that instead of just dropping the penny, they would introduce a 3 cent coin.
    - Like British three pence.
  - What is the largest amount that cannot be paid by using only existing coins (5, 10, 25) and a 3c coin?