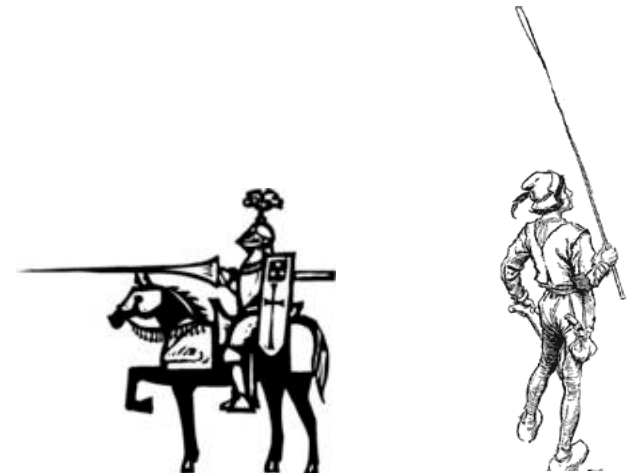


COMP 1002

Intro to Logic for Computer Scientists

Lecture 11





Puzzle 10



- The first formulation of the famous liar's paradox, attributed to a Cretan philosopher Epimenides, stated

“All Cretans are liars”.

Is this really a paradox?





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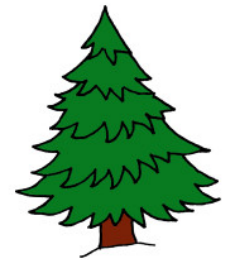
- *The negation of “all” is “exists”,*
 - *just like the negation of “and” is “or”*
- *So if Epimenides lied, what is true is that there are **some** truth-tellers on Crete (and potentially some liars, too)*
- *And Epimenides is one of the liars.*
- *However, “I am lying” would be a paradox.*



“NOT” makes life harder



- It is easy to visualize a tree, a number, or a person. It is harder to visualize a “not a tree”, “not a number” or “not a person”
- So “NOT (ALL trees have leaves)” is harder to understand than “some trees have something other than leaves (e.g., needles).”
- Here we really need to pay attention to the domain of quantifiers! It stays the same when negating.



– Not all integers are even: $\neg(\forall x \in \mathbb{Z} \text{ Even}(x))$

\equiv

– Some integers are not even $\exists x \in \mathbb{Z} \neg \text{Even}(x)$

Mixing quantifiers



- We can make statements of predicate logic mixing existential and universal quantifiers.
- Order of variables under the same quantifier does not matter. Under different ones does.

- Predicate: Loves(x,y). Domain: people.
- Everybody loves somebody: $\forall x \exists y \text{ Loves}(x,y)$
 - Normal people



Somebody loves everybody: $\exists x \forall y \text{ Loves}(x,y)$

- Mother Teresa

- Everybody is loved by somebody $\forall x \exists y \text{ Loves}(y,x)$
 - Their mother



- Somebody is loved by everybody $\exists x \forall y \text{ Loves}(y,x)$
 - Elvis Presley

- Everybody is loved by everybody $\forall x \forall y \text{ Loves}(x,y)$
 - Domain is a good family (not Meow-stery family)



Negating mixed quantifiers



- Now, a “not” in front of such a sentence means all \forall and \exists are interchanged, and the inner part becomes negated.

– Everybody loves somebody: $\forall x \exists y \text{ Loves}(x,y)$



- Somebody does not love anybody $\exists x \forall y \neg \text{Loves}(x,y)$
- Can also say “Somebody loves nobody” in English.
- Not the same as “somebody does not love everybody”:
here, “somebody does not (love everybody)” meaning
 $\exists x \neg (\forall y \text{ Loves}(x,y)) \equiv \exists x \exists y \neg \text{Loves}(x,y)$
- But the formula $\exists x \exists y \neg \text{Loves}(x,y)$ is the negation of
 $\forall x \forall y \text{ Loves}(x,y)$



Negating mixed quantifiers



- Everybody loves somebody: $\forall x \exists y \text{ Loves}(x,y)$
 - Somebody does not love anybody $\exists x \forall y \neg \text{Loves}(x,y)$



- Somebody loves everybody: $\exists x \forall y \text{ Loves}(x,y)$
 - Everyone doesn't like somebody $\forall x \exists y \neg \text{Loves}(x,y)$



- Everybody is loved by somebody $\forall x \exists y \text{ Loves}(y,x)$
 - Somebody is not loved by anybody $\exists x \forall y \neg \text{Loves}(y,x)$



- Somebody is loved by everybody $\exists x \forall y \text{ Loves}(y,x)$
 - For everyone, somebody does not love them $\forall x \exists y \neg \text{Loves}(y,x)$



- Everybody is loved by everybody $\forall x \forall y \text{ Loves}(y,x)$
 - Somebody does not love someone $\exists x \exists y \neg \text{Loves}(y,x)$



Scope of quantifiers



- Like in programming, a scope of a quantified variable continues until a new variable with the same name is introduced.
 - $\forall x (\exists y P(x, y)) \wedge (\exists y Q(x, y))$
 - For everybody there is somebody who loves them and somebody who hates them.
 - Not the same as $\forall x (\exists y P(x, y) \wedge Q(x, y))$
 - For everybody there is somebody who both loves and hates them.

- Better to avoid using same names for different variables – it is confusing.

$$\begin{aligned} & \text{– } \forall x (\exists y P(x, y)) \wedge (\exists y Q(x, y)) \\ & \quad \equiv \end{aligned}$$

$$\begin{aligned} & \text{– } \forall x (\exists y P(x, y)) \wedge (\exists z Q(x, z)) \\ & \quad \equiv \end{aligned}$$

$$\begin{aligned} & \text{– } \forall x \exists y \exists z P(x, y) \wedge Q(x, z) \\ & \quad \equiv \end{aligned}$$

$$\text{– } \forall x \exists z \exists y P(x, y) \wedge Q(x, z)$$

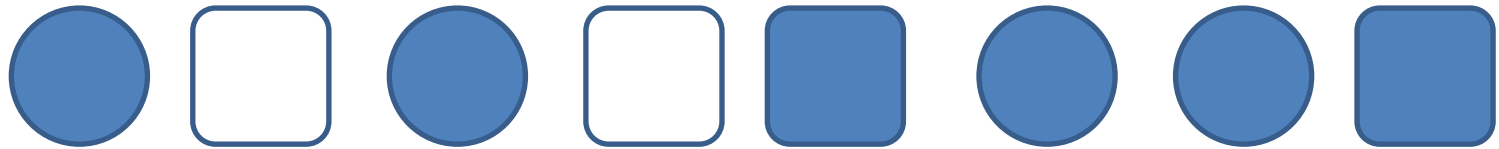
Equivalence for predicate logic



- Two predicate logic formulas are equivalent if they have the same truth value for every setting of free variables, no matter what the predicates are.
 - $(\exists y P(x, y)) \wedge (\exists y Q(x, y))$
 \equiv
 - $(\exists y P(x, y)) \wedge (\exists z Q(x, z))$
 \equiv
 - $\exists y \exists z P(x, y) \wedge Q(x, z)$
 \equiv
 - $\exists z \exists y P(x, y) \wedge Q(x, z)$
 - But $\exists x \forall y P(x, y, z)$ is not equivalent to $\forall y \exists x P(x, y, z)$

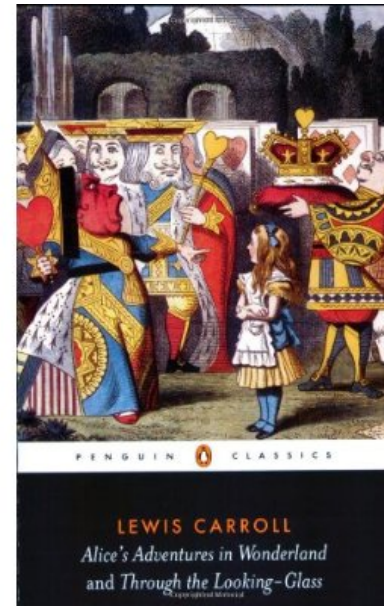
Quantifiers and conditionals

- Which statements are true?
 - All squares are white. All white shapes are squares
 - All circles are blue. All blue shapes are circles.



- All lemurs live in the trees. All animals living in the trees are lemurs.
- $\forall x \in S, P(x) \rightarrow Q(x)$
 - For all objects, if it is white, then it is a square.
 - If an object is white, then it is a square.
 - If an animal is a lemur, then it lives in the trees.

- Then you should say what you mean,' the March Hare went on.
- 'I do,' Alice hastily replied; 'at least—at least I mean what I say—that's the same thing, you know.'
- 'Not the same thing a bit!' said the Hatter. 'You might just as well say that "I see what I eat" is the same thing as "I eat what I see"!'
- 'You might just as well say,' added the March Hare, 'that "I like what I get" is the same thing as "I get what I like"!'
- 'You might just as well say,' added the Dormouse, who seemed to be talking in his sleep, 'that "I breathe when I sleep" is the same thing as "I sleep when I breathe"!'



"Alice's Adventures in Wonderland"
by Lewis Carroll