Computer Science 3600 (Winter 2024):
Assignment #3
Supplementary Question Answers

2. (15 marks) Determine the optimal parenthesization of a matrix-chain product whose sequence of dimensions is $\langle 2, 5, 3, 4, 2, 3, 5\rangle$ using the algorithm given on page 336 of the textbook. In the style of Figure 15.3, show the filled-in dynamic programming matrices m and s , the "backpointer path" in s that gives an optimal parenthesization, and the parenthesization associated with that path.

Answer: The matrices m and s are shown in Figure 1. For the sake of completeness, the computations for each of the cells in m and s are given below:

$$m[1][2] = \min \left\{ \begin{array}{ll} m[1][1] + m[2][2] + (p[0] * p[1] * p[2]) & (k = 1) \\ = 0 + 0 + 30 = 30 & \end{array} \right\} = 30 \ (k = 1)$$

$$m[2][3] = \min \left\{ \begin{array}{ll} m[2][2] + m[3][3] + (p[1] * p[2] * p[3]) & (k = 2) \\ = 0 + 0 + 60 = 60 & \end{array} \right\} = 60 \ (k = 2)$$

$$m[3][4] = \min \left\{ \begin{array}{ll} m[3][3] + m[4][4] + (p[2] * p[3] * p[4]) & (k = 3) \\ = 0 + 0 + 24 = 24 & \end{array} \right\} = 24 \ (k = 3)$$

$$m[4][5] = \min \left\{ \begin{array}{ll} m[4][4] + m[5][5] + (p[3] * p[4] * p[5]) & (k = 4) \\ = 0 + 0 + 24 = 24 & \end{array} \right\} = 24 \ (k = 4)$$

$$m[5][6] = \min \left\{ \begin{array}{ll} m[5][5] + m[6][6] + (p[4] * p[5] * p[6]) & (k = 5) \\ = 0 + 0 + 30 = 30 & \end{array} \right\} = 30 \ (k = 5)$$

$$m[1][3] = \min \left\{ \begin{array}{ll} m[1][1] + m[2][3] + (p[0] * p[1] * p[3]) & (k = 1) \\ = 0 + 60 + 40 = 100, & \\ m[1][2] + m[3][3] + (p[0] * p[2] * p[3]) & (k = 2) \\ = 30 + 0 + 24 = 54 & \end{array} \right\} = 54 \ (k = 2)$$

$$m[2][4] = \min \left\{ \begin{array}{ll} m[2][2] + m[3][4] + (p[1] * p[2] * p[4]) & (k = 2) \\ = 0 + 24 + 30 = 54, & \\ m[2][3] + m[4][4] + (p[1] * p[3] * p[4]) & (k = 3) \\ = 60 + 0 + 40 = 100 & \end{array} \right\} = 54 \ (k = 2)$$

$$m[3][5] = \min \left\{ \begin{array}{ll} m[3][3] + m[4][5] + (p[2] * p[3] * p[5]) & (k = 3) \\ = 0 + 24 + 36 = 60, & \\ m[3][4] + m[5][5] + (p[2] * p[4] * p[5]) & (k = 4) \\ = 24 + 0 + 18 = 42 & \end{array} \right\} = 42 \ (k = 4)$$

$$m[4][6] = \min \left\{ \begin{array}{ll} m[4][4] + m[5][6] + (p[3] * p[4] * p[6]) & (k = 4) \\ = 0 + 30 + 40 = 70, & \\ m[4][5] + m[6][6] + (p[3] * p[5] * p[6]) & (k = 5) \\ = 24 + 0 + 60 = 84 & \end{array} \right\} = 70 \ (k = 4)$$

$$m[1][4] = \min \left\{ \begin{array}{ll} m[1][1] + m[2][4] + (p[0] * p[1] * p[4]) & (k = 1) \\ = 0 + 54 + 20 = 74, & \\ m[1][2] + m[3][4] + (p[0] * p[2] * p[4]) & (k = 2) \\ = 30 + 24 + 12 = 66, & \\ m[1][3] + m[4][4] + (p[0] * p[3] * p[4]) & (k = 3) \\ = 54 + 0 + 16 = 70 & \end{array} \right\} = 66 \ (k = 2)$$

$$m[2][5] = \min \left\{ \begin{array}{ll} m[2][2] + m[3][5] + (p[1] * p[2] * p[5]) & (k = 2) \\ = 0 + 42 + 45 = 87, & \\ m[2][3] + m[4][5] + (p[1] * p[3] * p[5]) & (k = 3) \\ = 60 + 24 + 60 = 144, & \\ m[2][4] + m[5][5] + (p[1] * p[4] * p[5]) & (k = 4) \\ = 54 + 0 + 30 = 84 & \end{array} \right\} = 84 \ (k = 4)$$

$$m[3][6] = \min \left\{ \begin{array}{ll} m[3][3] + m[4][6] + (p[2] * p[3] * p[6]) & (k = 3) \\ = 0 + 70 + 60 = 130, & \\ m[3][4] + m[5][6] + (p[2] * p[4] * p[6]) & (k = 4) \\ = 24 + 30 + 30 = 84, & \\ m[3][5] + m[6][6] + (p[2] * p[5] * p[6]) & (k = 5) \\ = 42 + 0 + 45 = 87 & \end{array} \right\} = 84 \ (k = 4)$$

$$m[1][5] = \min \left\{ \begin{array}{ll} m[1][1] + m[2][5] + (p[0] * p[1] * p[5]) & (k = 1) \\ = 0 + 84 + 30 = 114, & \\ m[1][2] + m[3][5] + (p[0] * p[2] * p[5]) & (k = 2) \\ = 30 + 42 + 18 = 90, & \\ m[1][3] + m[4][5] + (p[0] * p[3] * p[5]) & (k = 3) \\ = 54 + 24 + 24 = 102, & \\ m[1][4] + m[5][5] + (p[0] * p[4] * p[5]) & (k = 4) \\ = 66 + 0 + 12 = 78 & \end{array} \right\} = 78 \ (k = 4)$$

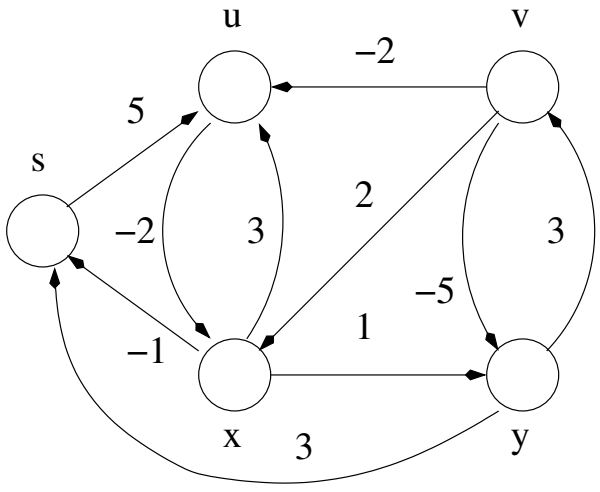
$$m[2][6] = \min \left\{ \begin{array}{ll} m[2][2] + m[3][6] + (p[1] * p[2] * p[6]) & (k = 2) \\ = 0 + 84 + 75 = 159, & \\ m[2][3] + m[4][6] + (p[1] * p[3] * p[6]) & (k = 3) \\ = 60 + 70 + 100 = 230, & \\ m[2][4] + m[5][6] + (p[1] * p[4] * p[6]) & (k = 4) \\ = 54 + 30 + 50 = 134, & \\ m[2][5] + m[6][6] + (p[1] * p[5] * p[6]) & (k = 5) \\ = 84 + 0 + 75 = 159 & \end{array} \right\} = 134 \ (k = 4)$$

$$m[1][6] = \min \left\{ \begin{array}{ll} m[1][1] + m[2][6] + (p[0] * p[1] * p[6]) & (k = 1) \\ = 0 + 134 + 50 = 184, & \\ m[1][2] + m[3][6] + (p[0] * p[2] * p[6]) & (k = 2) \\ = 30 + 84 + 30 = 144, & \\ m[1][3] + m[4][6] + (p[0] * p[3] * p[6]) & (k = 3) \\ = 54 + 70 + 40 = 164, & \\ m[1][4] + m[5][6] + (p[0] * p[4] * p[6]) & (k = 4) \\ = 66 + 30 + 20 = 116, & \\ m[1][5] + m[6][6] + (p[0] * p[5] * p[6]) & (k = 5) \\ = 78 + 0 + 30 = 108 & \end{array} \right\} = 108 \ (k = 5)$$

The "backpointer path" (more properly, the backpointer tree) is indicated by dashed arrows linking cells in the s matrix in Figure 1. This tree implies that the optimal parenthesization of the given matrix chain is

$$\begin{array}{l} (A_1A_2A_3A_4A_5A_6) \\ \Rightarrow ((A_1A_2A_3A_4A_5)(A_6))) \\ \Rightarrow (((A_1A_2A_3A_4)(A_5))(A_6))) \\ \Rightarrow ((((A_1A_2)(A_3A_4))(A_5))(A_6))) \end{array}$$

3. (20 marks) Consider the following edge-weighted directed graph:



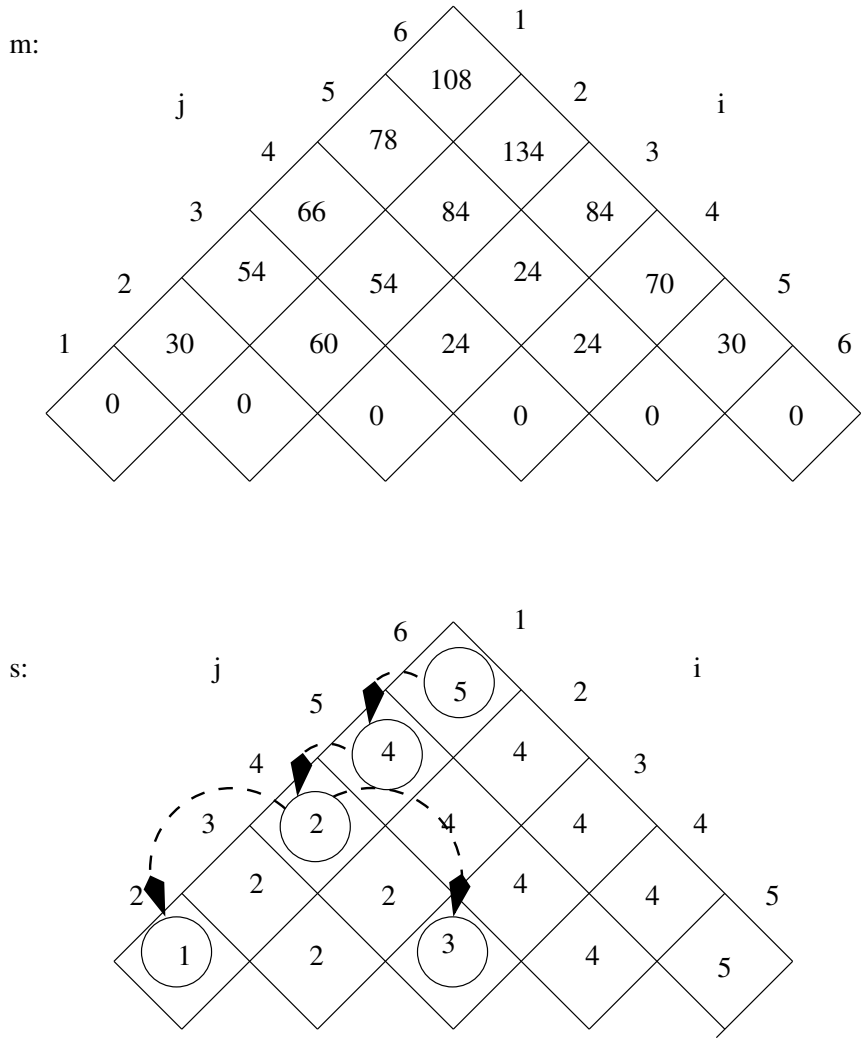


Figure 1: Answer for Question #2.

a) (10 marks) Run Dijkstra's algorithm (p, 595) on the directed graph above using vertex y as the source vertex. In the style of Figure 24.6 in the textbook, show the d and π values and the vertices in set S after each iteration of the **while** loop.

Answer: See Table 1.

b) (10 marks) Run the Bellman-Ford algorithm (p, 588) on the directed graph above using vertex y as the source vertex. Relax edges in lexicographic order in each pass, and in the style of Figure 24.4 on the textbook, show the d and π values after each pass. Finally, give the boolean value returned by the algorithm.

Answer: See Table 2. As it is the case for all edges (a, b) that $d(b) \leq w(a, b) + d(a)$, the algorithm returns `true`.

4. (20 marks) Consider the following decision problems:

Dumbbell subgraph (DS)
Input: An undirected graph $G = (V, E)$ and two positive integers $k, l \geq 1$.
Question: Are there two cliques C_1 and C_2 and a simple path P in G such that C_1 and C_2 have $\geq k$ vertices apiece, P has $\geq l$ edges, P connects C_1 and C_2, the cliques and path do not have any edges in common, and the only vertices that P shares with C_1 (C_2) is its connection-vertex?

Bounded-Weight Subset Cover (BWSSC)
Input: A set $I = \{i_1, \ldots, i_a\}$ of items, a set $R = \{r_1, \ldots, r_b\}$ of subsets of I, an integer-valued subset-weight function $w()$ such that for each $r_x \in R$, $w(r_x) > 0$, a subset $N \subseteq I$, and integers $0 < k_1 \leq k_2$.
Question: Is there is a subset $R' \subseteq R$ such that $\cup_{r \in R'} r = N$ and $k_1 \leq \sum_{r \in R'} w(r) \leq k_2$?

a) (10 marks) Prove that problem DS is NP-complete by (1) showing that this problem is in NP and (2) giving a polynomial-time many-one reduction (algorithm + proof of correctness) to this problem from an NP-hard problem.

Answer: As any candidate solution has size s such that $\frac{2k(k-1)}{2} \leq s \leq |V|$ and this solution can be checked in time polynomial in the input size (ensure all vertices are distinct and that the required edges exist between them), this problem is in NP. To show NP-hardness, we reduce from the following NP-complete problem (p. 983, textbook):

Hamiltonian Path (HP)
Input: An undirected graph G.

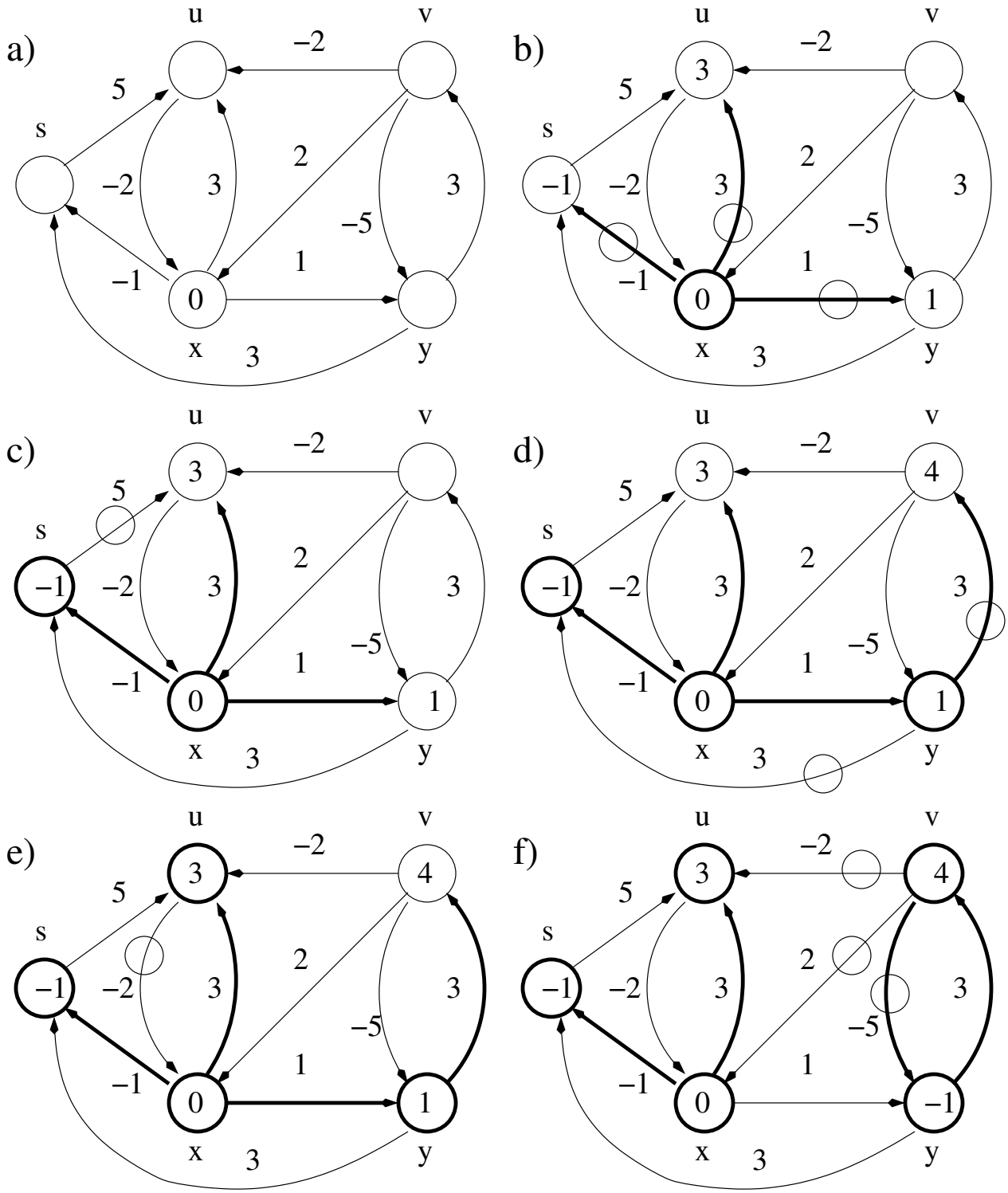


Table 1: Answer for Question #3(a). The shortest-path estimates are shown within the vertices; empty vertices have estimates equal to infinity (∞). Bold edges indicate predecessor values. Bold vertices are in the set S and regular vertices are in the priority queue $Q = V - S$. **(a)** The situation before the first execution of the **while** loop on lines 4–8. **(b)**–**(f)** The situation after each successive iteration of the **while** loop. Note that edges for which relaxation is attempted in each successive iteration of the **while** loop are marked with circles.

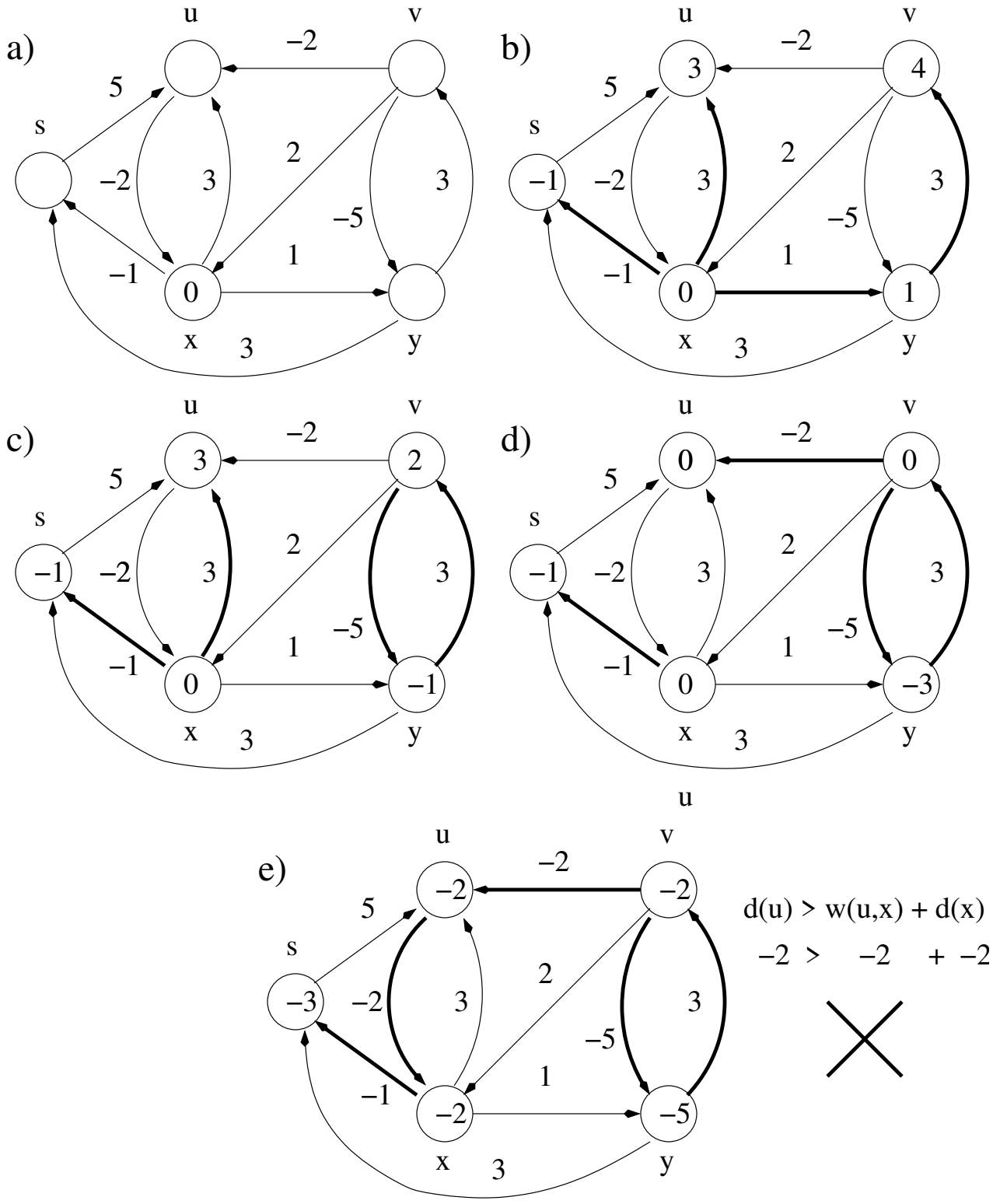


Table 2: Answer for Question #3(b). The shortest-path estimates are shown within the vertices; empty vertices have estimates equal to infinity (∞). Bold edges indicate predecessor values. **(a)** The situation just before the first pass over the edges. **(b)**–**(e)** The situation after each successive pass over the edges.

Question: Is there a simple path in G that links all vertices in G?

Given and instance (G) of HP, we construct an instance (G, k, l) of DS in which $k = 1$ and $l = |V| - 1$. Any solution to this instance of DS contains a simple path with at least $|V| - 1$ edges which must by definition include all vertices in G and hence is a Hamiltonian path; conversely, any solution to the given instance of HP is a simple path in G with $|V| - 1$ edges which connects two cliques of size $k = 1$, *i.e.*, the endpoints of the path. As the construction described above can also be done in polynomial time, this construction is a polynomial-time many-one reduction from HP to DS which establishes the NP-hardness of DS. As DS is NP-hard and also in NP, DS is NP-complete.

b) (10 marks) Prove that problem BWSSC is NP-complete by (1) showing that this problem is in NP and (2) giving a polynomial-time many-one reduction (algorithm + proof of correctness) to this problem from an NP-hard problem.

Answer: As any candidate solution for an instance of BWSSC can be checked in time polynomial in the input size (ensure that the selected subset of subsets covers N and that the sum of the weights of these selected subsets is between k_1 and k_2 inclusive), this problem is in NP. To show NP-hardness, we reduce from the following NP-complete problem (p. 1033–1034, textbook):

Set Cover (SC)
Input: A set $U = \{u_1, \ldots, u_n\}$ of items, a set $S = \{s_1, \ldots, s_m\}$ of subsets of U, and an integer $k \leq m$.
Question: Is there a subset $S' \subseteq S$ such that $|S| \leq k$ and $\cup_{s \in S'} s = U$?

Given an instance $\langle U, S, k \rangle$ of SC, we construct an instance $\langle I, R, w, N, k_1, k_2 \rangle$ of BWSSC such that $I = N = U$, $R = S$, $w(r) = 1$ for all $r \in R$, $k_1 = 0$, and $k_2 = k$. Observe that any solution to the constructed instance of BWSSC is also a solution to the given instance of SC, and vice versa. As the construction described above can be done in polynomial time, this construction is a polynomial-time many-one reduction from SC to BWSSC which establishes the NP-hardness of BWSSC. As BWSSC is NP-hard and also in NP, BWSSC is NP-complete.