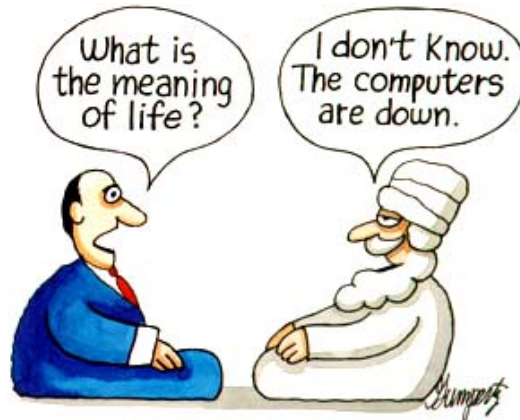
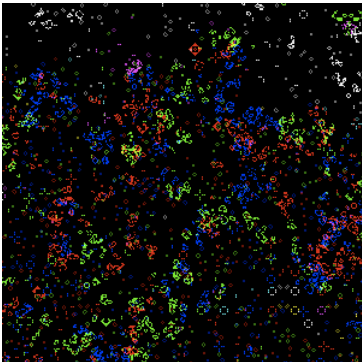


# Limits of Computation

Antonina Kolokolova

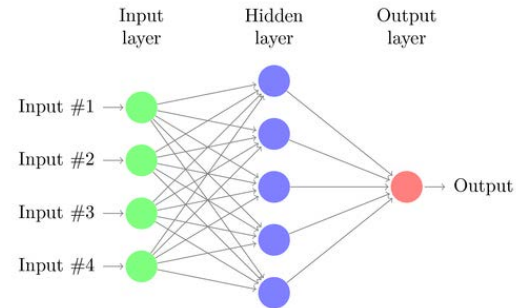


- What is computation?



- What is information?

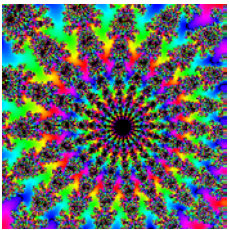
- What is learning?



- Are there any limits of our ability to solve problems?

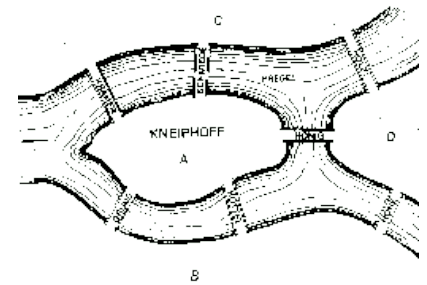
# Theoretical Computer Science

- Is there a perfect antivirus?



- Can computers be creative?

- Why some problems are easier than others?



- Is it possible to have secure information and communication?





# What is information?

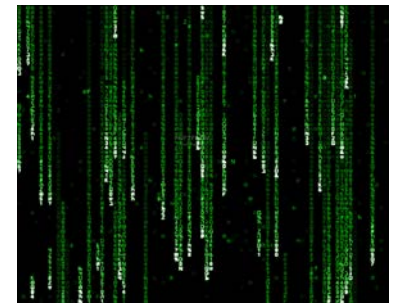


- Does string 1111111111 contain more information than the string 10010110100?
- Do you learn more from a coin toss of a fair coin or a roll of dice?





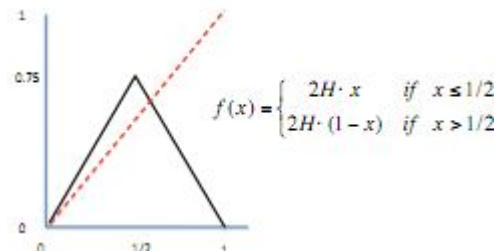
# What is information?



The less you can predict an outcome

The more you learn from it:

The more information you get.



# The science of information

- In many languages the word for “Computer Science” is derived from the word for information

- French: Informatique
- Spanish: Informática
- German: Informatik
- Russian: Информатика

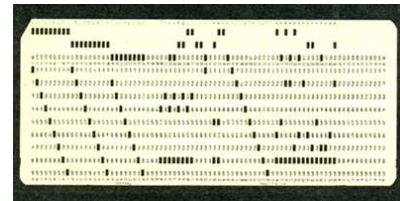


- The information comes in and we process it.
- So do computers. So do living cells, etc, etc.



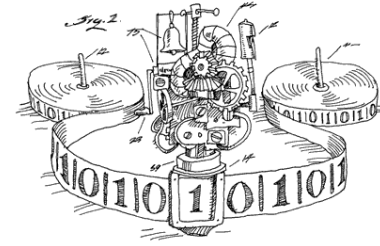
# On the other side of iron curtain

- In Soviet Union, in particular in Ukraine, PCs were not around till 1990<sup>th</sup>
- First photo: “MIR” computer (from 1969). Developed in Kiev by Glushkov and his group.
- Were still in use in 1980s.
- Programmable calculators for personal use





# What is computation?



- We process information by doing a “computation on it”. Changing it from one representation to another.

- But what is computation?

- What does your smartphone compute when you are playing Pokemon Go?

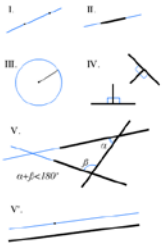


- How does DNA “compute”?

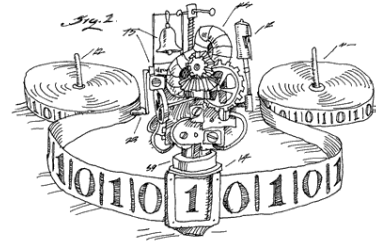


- Is there a limit to what can be computed?



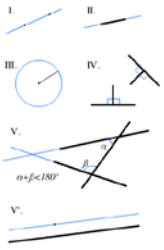


# Limits of computation

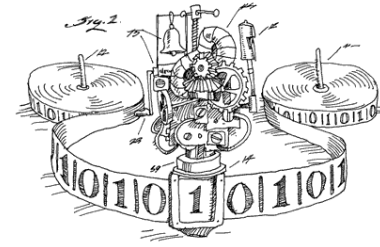


- In 1900, at the International Congress of Mathematicians in Paris, David Hilbert posed a list of 23 problems. Problem 2 asked to prove that mathematics contains no self-contradictions.
- In 1920, Hilbert extended it to what is now known as “Hilbert’s program”

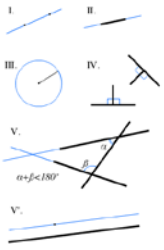




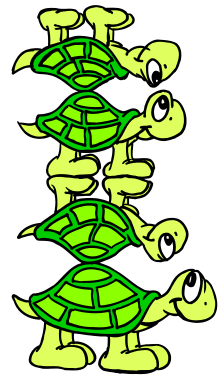
# Hilbert's program



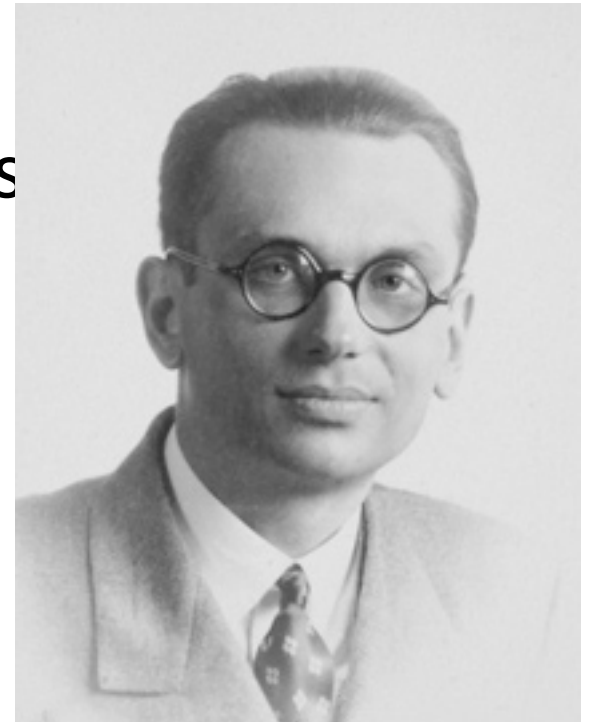
- Express all mathematics in a precise way
- Allowing a formal proof of all true statements
- With a proof, inside mathematics, that there is no self-contradiction
- And a procedure (an algorithm) for deciding, for any given mathematical statement, whether it is true or false.

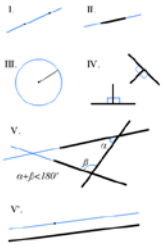


# Gödel Incompleteness Theorem

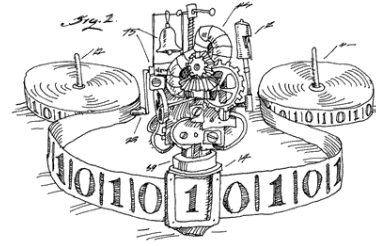


- If mathematics is not self-contradictory...
- Then there are true statements that can't be proven!
- Such as “I am not provable”
- A paradox!





# Church and Turing:

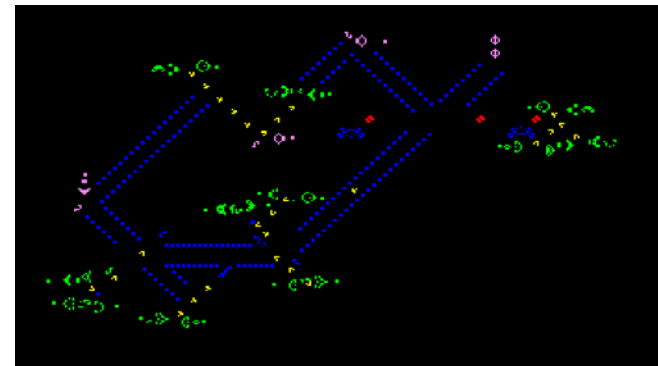
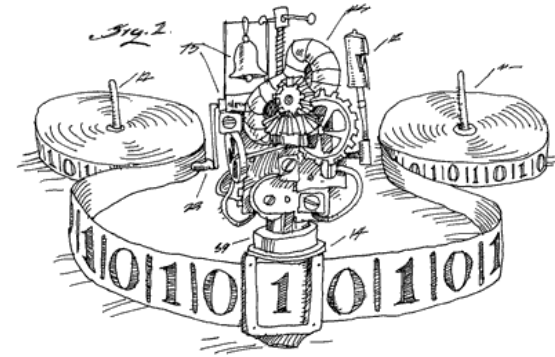


- Moreover,
- there is no procedure
- to decide if a given statement is true or false!
- And to decide many other things...
  
- But what do we mean by a “procedure”?



# Models of computation

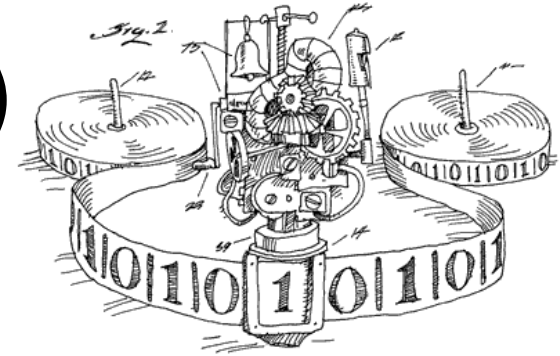
- Let me show you two models of computation.
- The first one is the Turing machine
  - Our modern-day computers are based on this model
- The second is the Game of Life
  - Looks nothing like a computer, and yet has the same power.





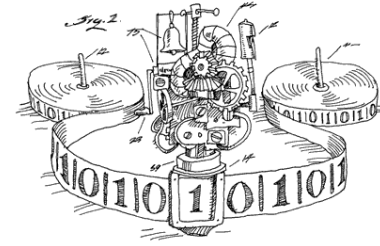
# Turing machine

- A Turing machine has an (unlimited) memory, visualized as a tape
- Or a stack of paper
- And takes very simple instructions:
  - Read a symbol
  - Write a symbol
  - Move one step left or right on the tape
  - Change internal state.

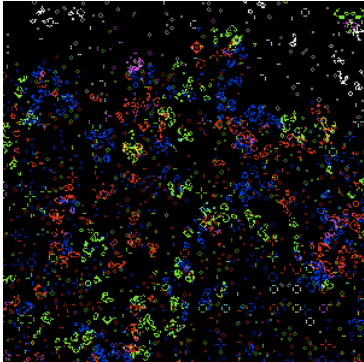




# Church-Turing thesis

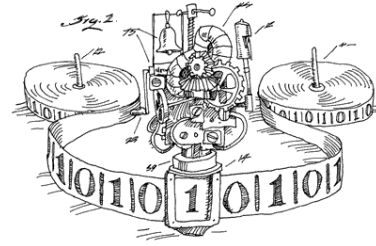


***Everything we can call “computable” is computable by a Turing machine.***

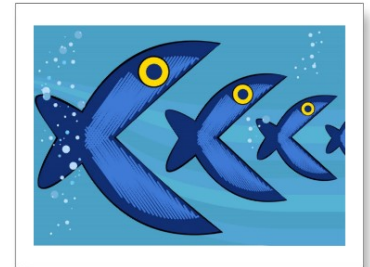




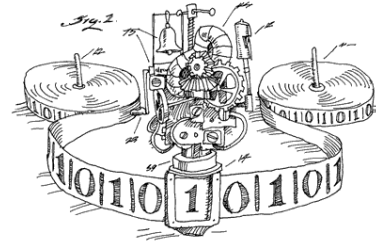
# “Will this ever stop?”



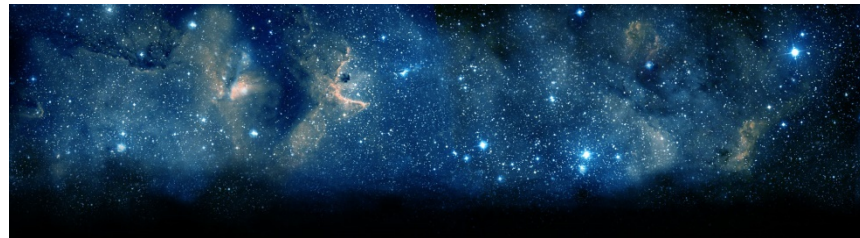
- A code for a Turing machine program is a string.
- Any string can be an input to a program.
- Imagine there is a machine that always does the opposite...
  - From the machine which code is its input
    - On a string encoding it
- What will it do on its own code?
  - Yes?... No?... Yes?... No?... Paradox!
- So no such machine can exist... Some problems are unsolvable, with self-reference to blame.



# Complexity of computation

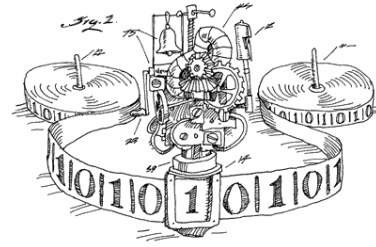


- Would you still consider a problem really solvable if it takes very long time?
  - Say  $10^n$  steps on an  $n$ -symbol string?
  - At a billion ( $10^9$ ) steps per second ( $\sim 1$ GHz)?
  - To process a string of length 100...
  - will take  $10^{100}/10^9$  seconds, or  $\sim 3 \times 10^{72}$  centuries.



- Age of the universe: about  $1.38 \times 10^{10}$  years.
- Atoms in the observable universe:  $10^{78}$ - $10^{82}$ .

# Complexity of computation



- What strings do we work with in real life?

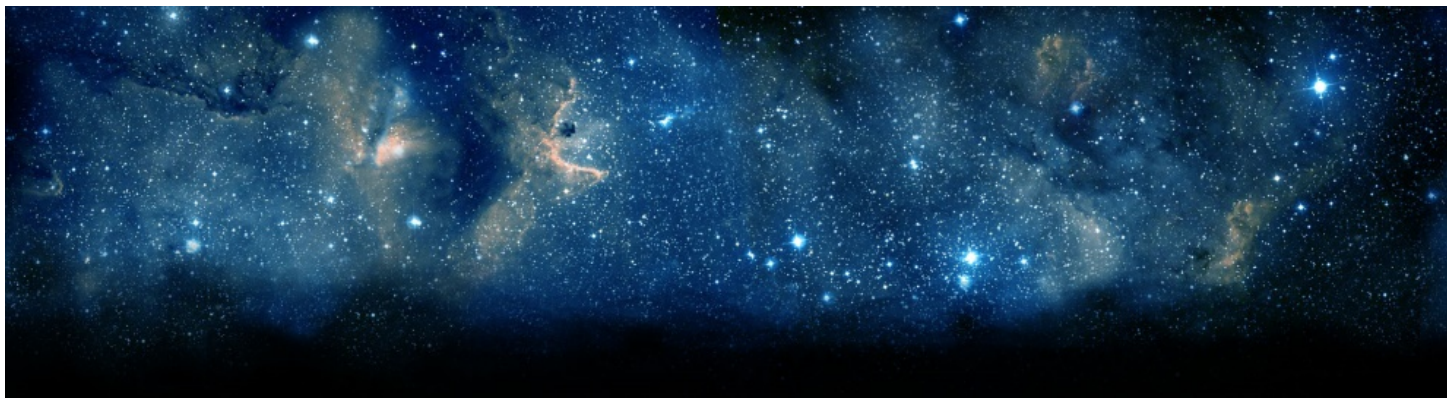
- A DNA string has  $3.2 \times 10^9$  base pairs



- A secure key in crypto: 128-256 bits

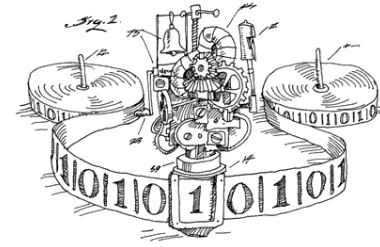
- Number of Walmart transactions per day:  $10^6$ .


- URLs searched by Google in 2012:  $3 \times 10^{12}$ .





# Efficient computation



- What can be computed in our universe?
  - We could only work with very short strings...
    - But we want to work with our DNA string! 
  - We can try being efficient in solving problems.
    - What does it mean to be efficient?
    - And what kinds of problems can be solved efficiently?

A million-dollar question!



# The million dollar question



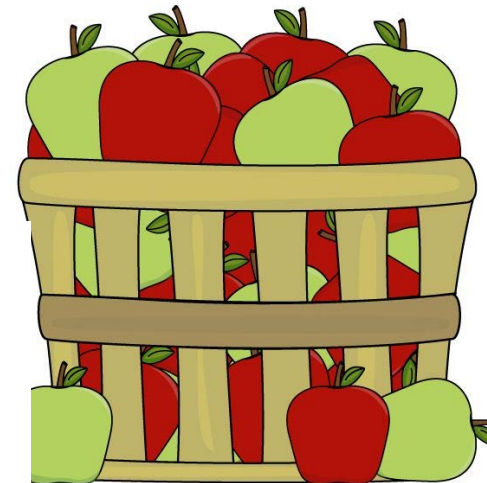
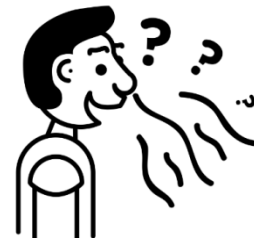
- In Russian, called “perebor” problem.
  - “perebor” translates as “exhaustive search”.
  - Question: is it always possible to avoid looking through nearly all potential solutions to find an answer?
    - Combinations of letters to guess your password?
    - Combinations of numbers to solve an equation?

**Are there situations when exhaustive search is unavoidable?**



# The million dollar question

- Suppose you have a basket of apples.
- Can you check that all apples are good without looking at (essentially) every single one?
  - Is there a way that would work for every possible basket of apples?
- **Smell test?**





# The million dollar question



- In English, most known as “P vs. NP” problem
  - P stands for “polynomial time computable”.
  - NP is “polynomial time checkable”
    - non-deterministic polynomial-time computable
    - (fancy word for “guess and check”)
- **Question: is everything efficiently checkable also efficiently computable?**
  - In particular, is there a “smell test” for every problem that has easy-to-check solutions?

# Colouring maps

- How many colours needed so that neighbouring countries do not get the same colour?
- For a picture like that – no more than 4.
- (A theorem famous for being proven with a help of a computer)
- Can it be done with 3?





# Colouring maps

- Can it be done with 3 colours?
- How do we find out?
  - Look at neighbours of Austria. There are 7 of them... 3 colours not enough.
- In general, nobody knows a good way!



# Question

- Can this map be coloured with three colours?



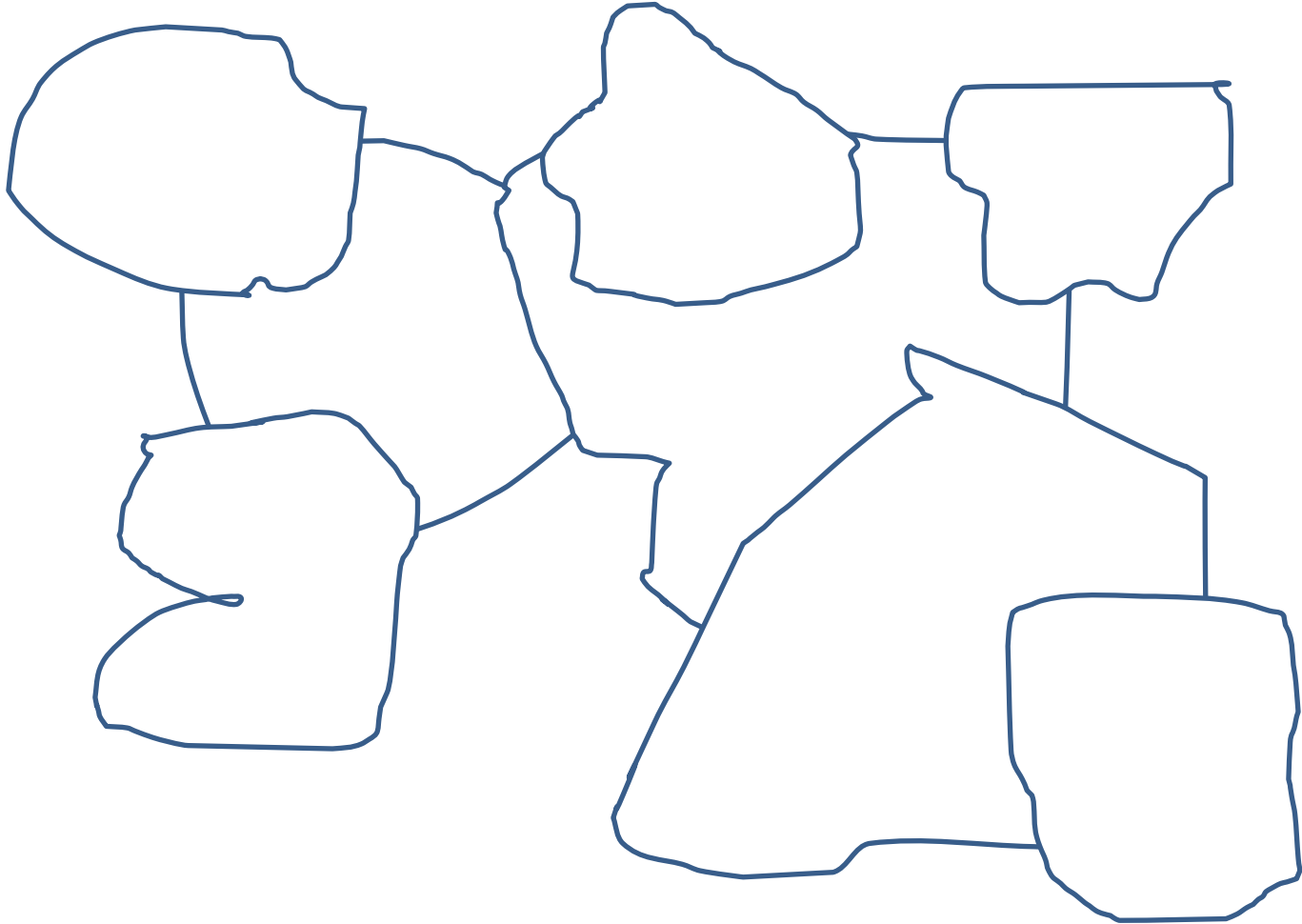
# Question

- Can this map be coloured with three colours?



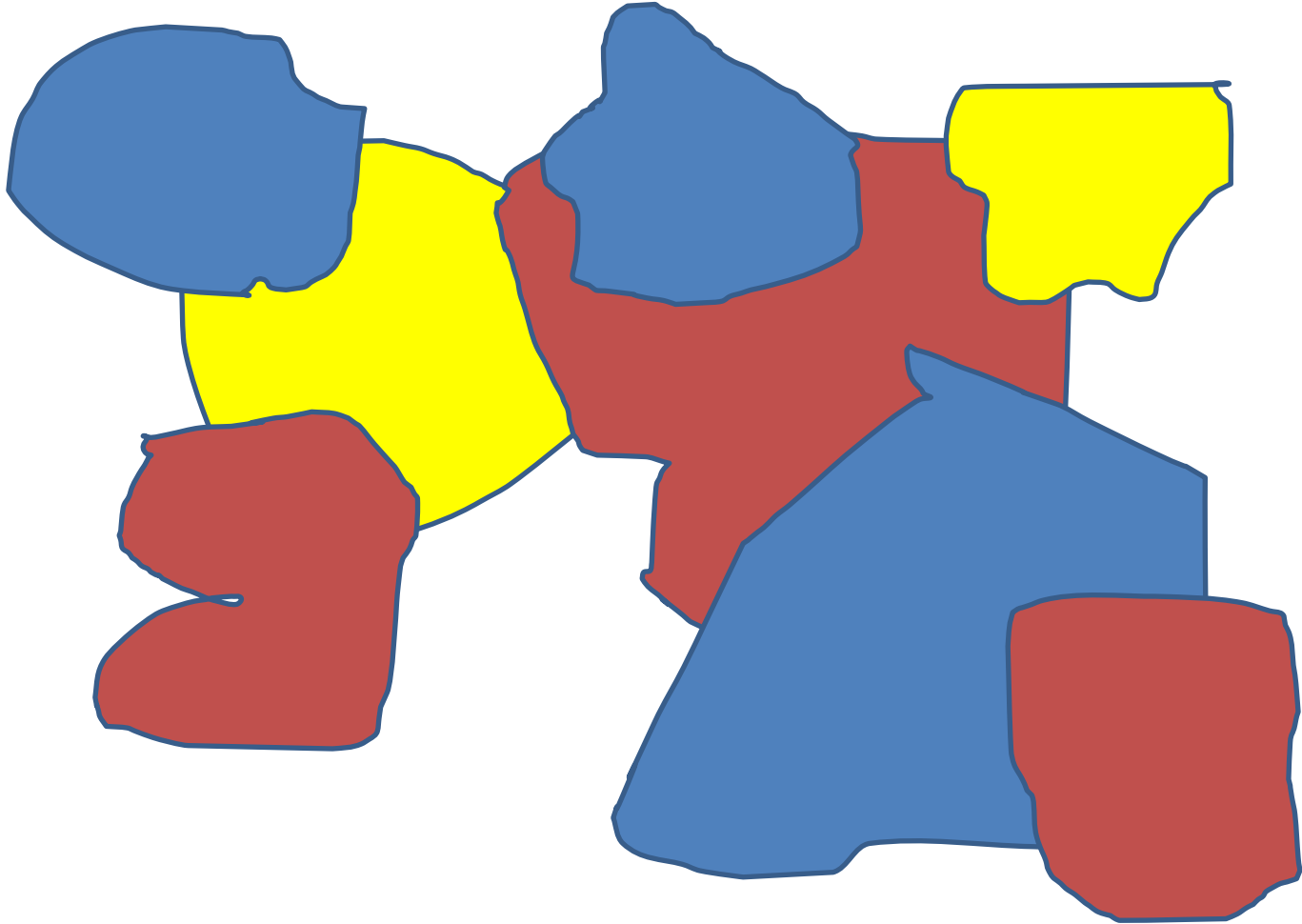
# Question

- Can this map be coloured with only 3 colours?



# Question

- Can this map be coloured with only 3 colours?





# Question

- Can this map be coloured with two colours?



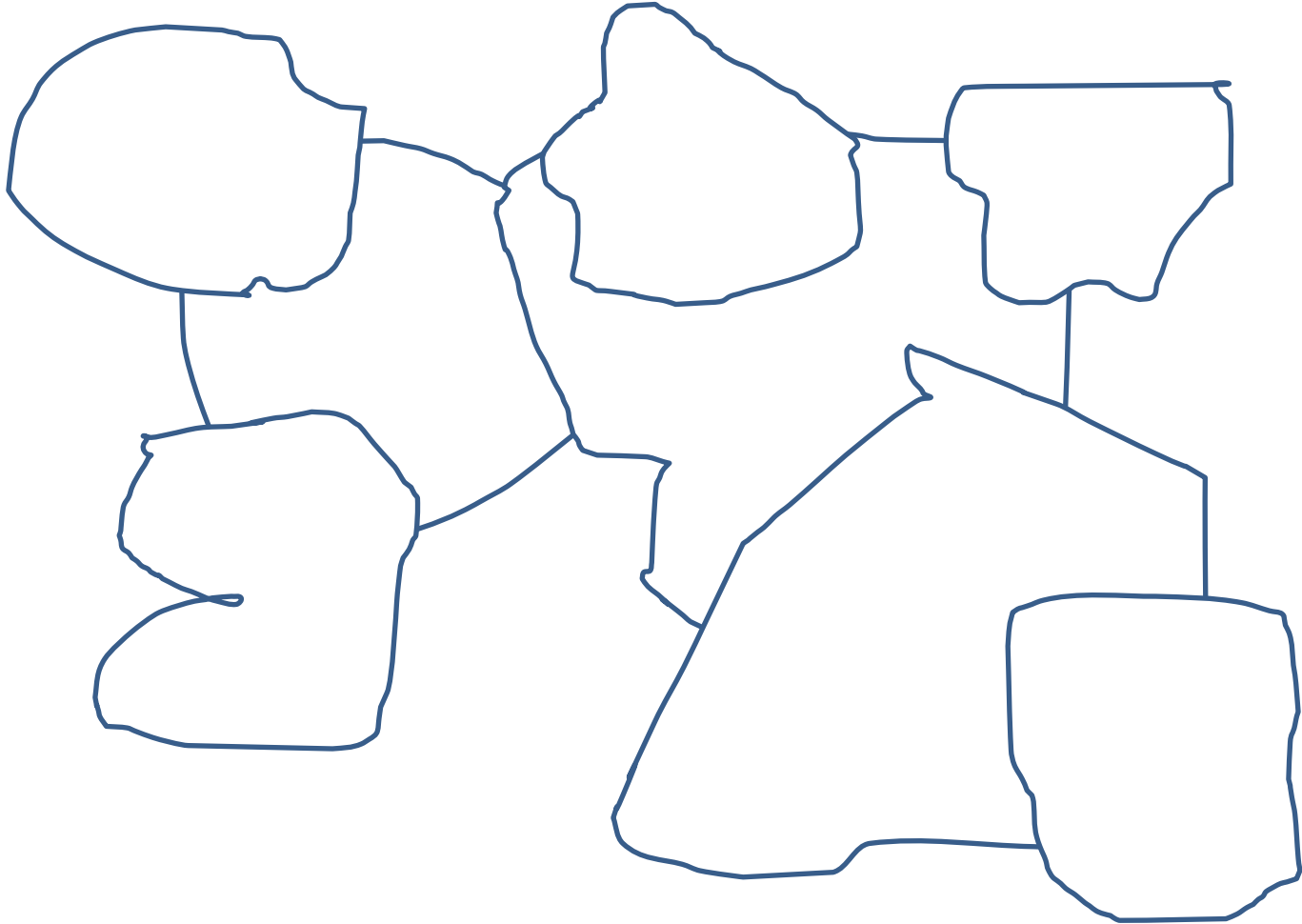
# Question

- Can this map be coloured with two colours?
- No...
- Western Australia, Northern territory and South Australia should all be different colours.



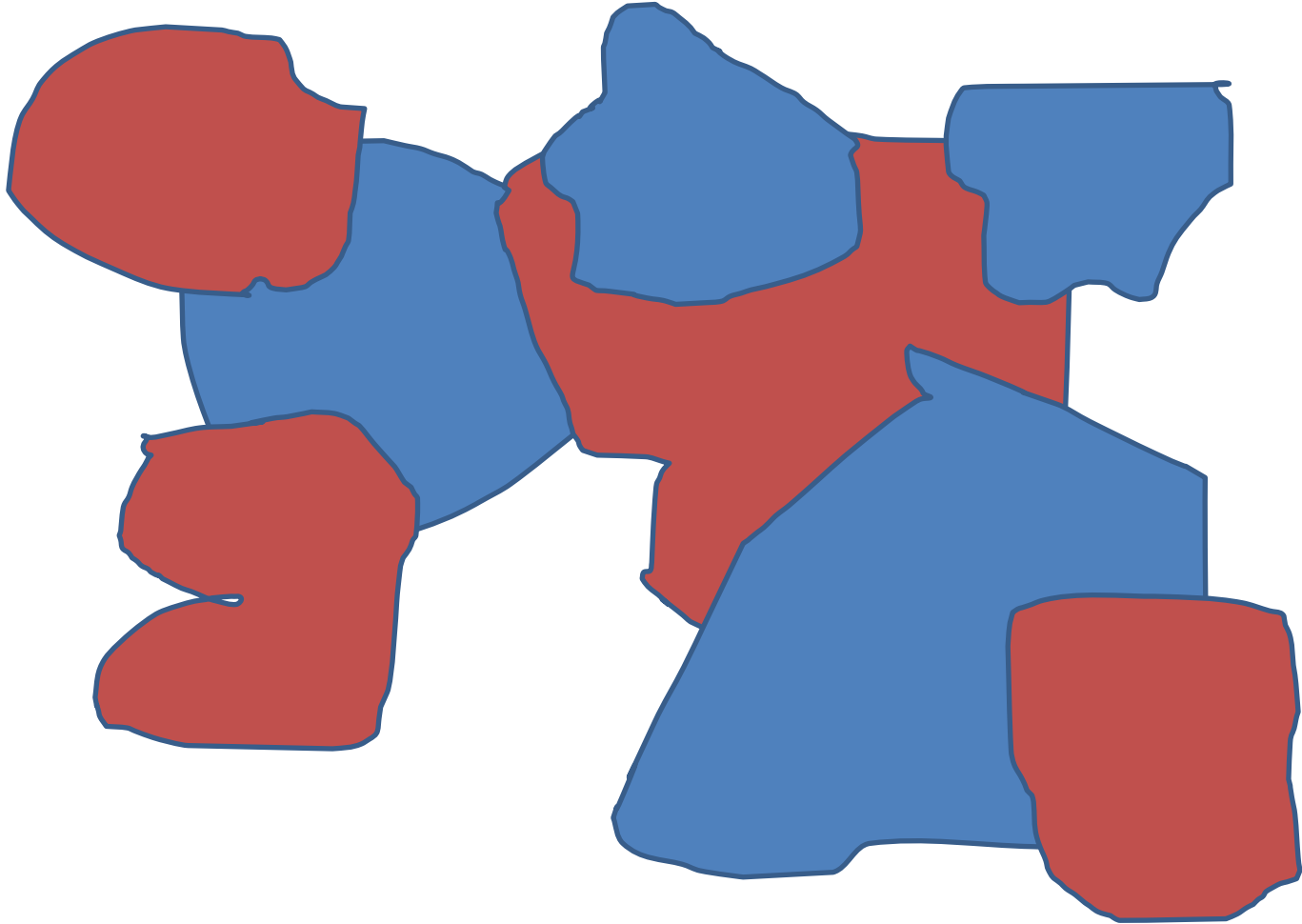
# Question

- Can this map be coloured with only 2 colours?



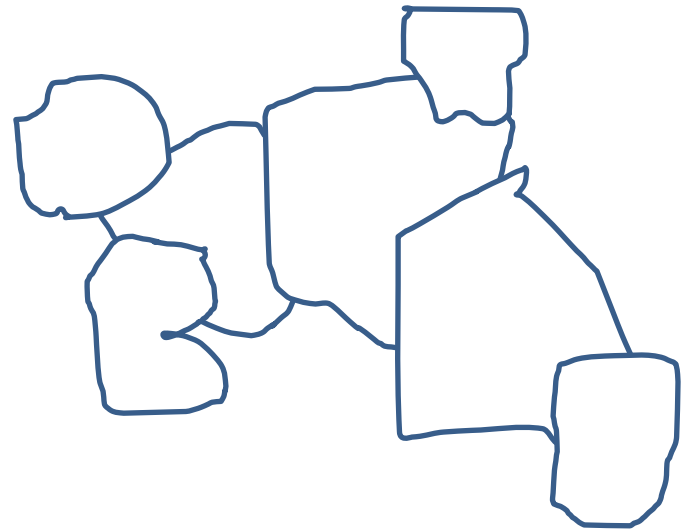
# Question

- Can this map be coloured with only 2 colours?



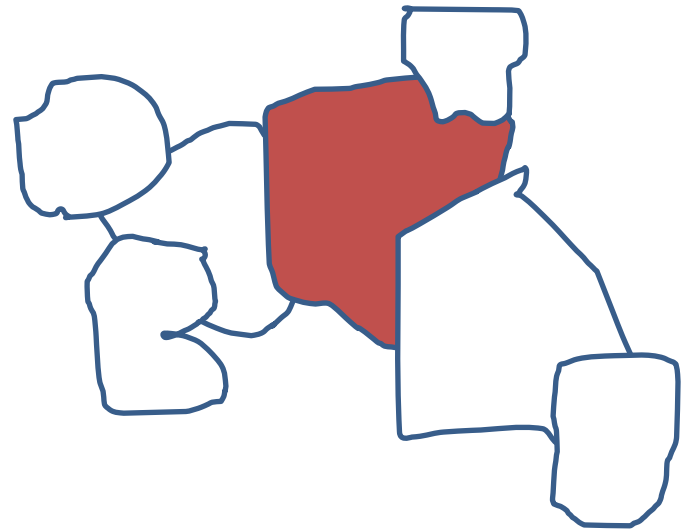
# Colouring with 2 colours

- How do you check if a map is colourable with 2 colours?
  - Start anywhere.
  - Colour a region red
  - Colour its neighbours blue
  - Colour their neighbours red again...
  - Continue until either done, or found a region would border one of the same colour



# Colouring with 2 colours

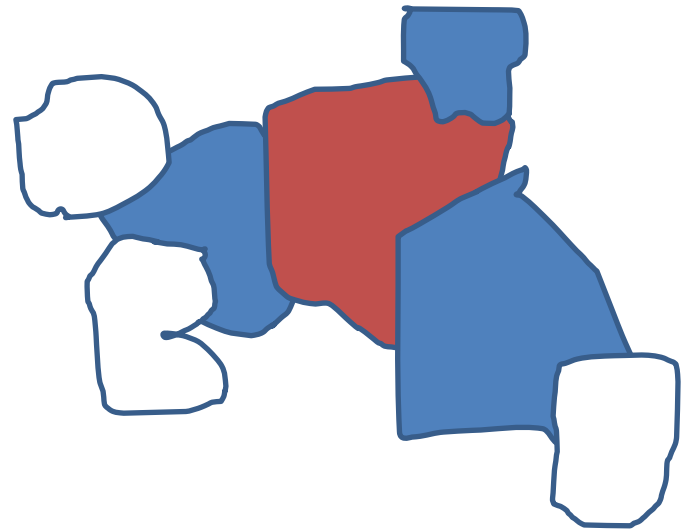
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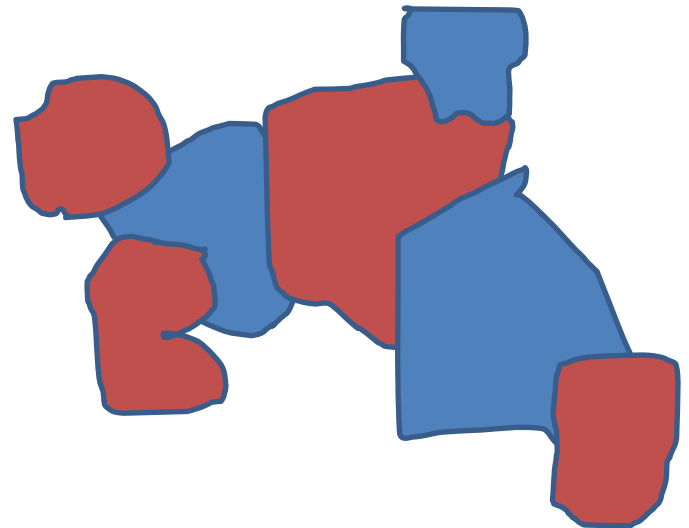
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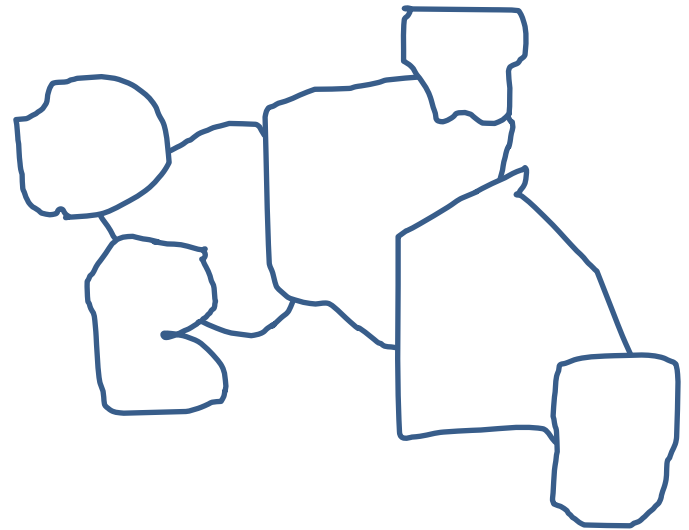
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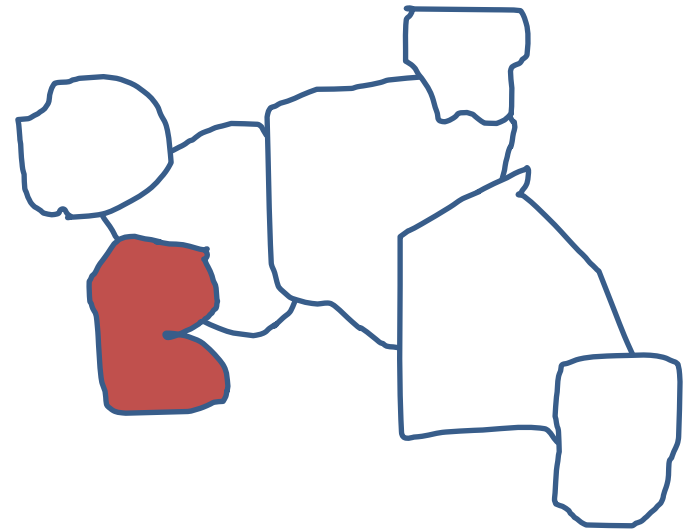
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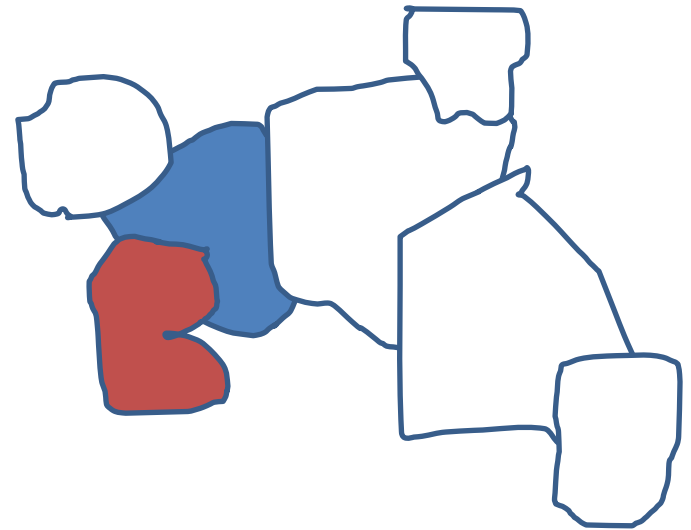
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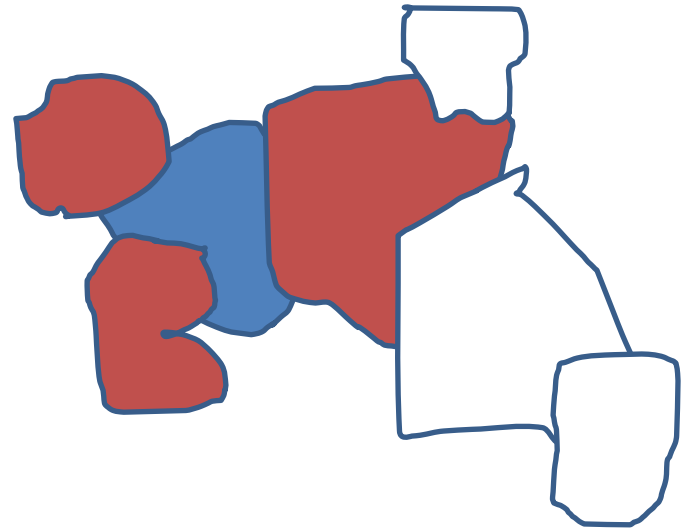
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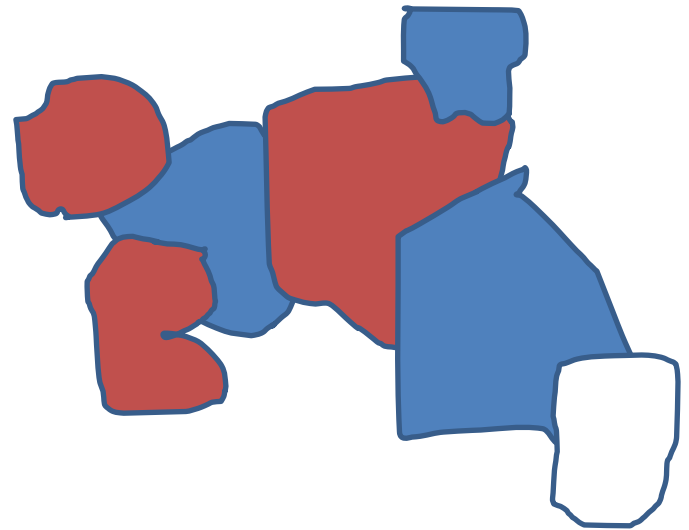
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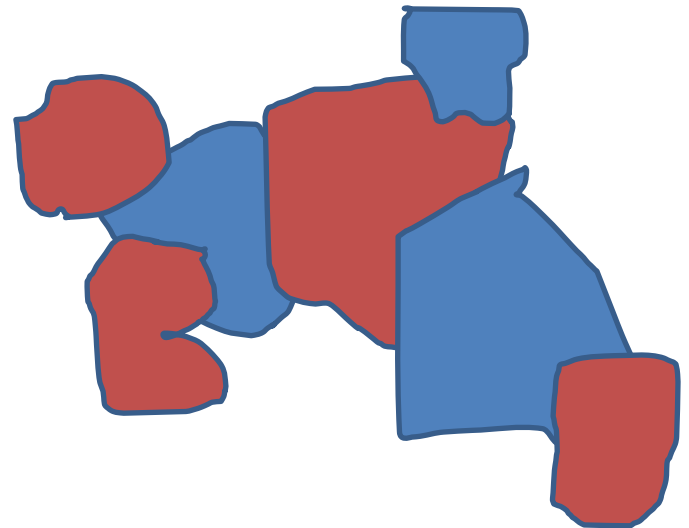
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# Colouring with 2 colours

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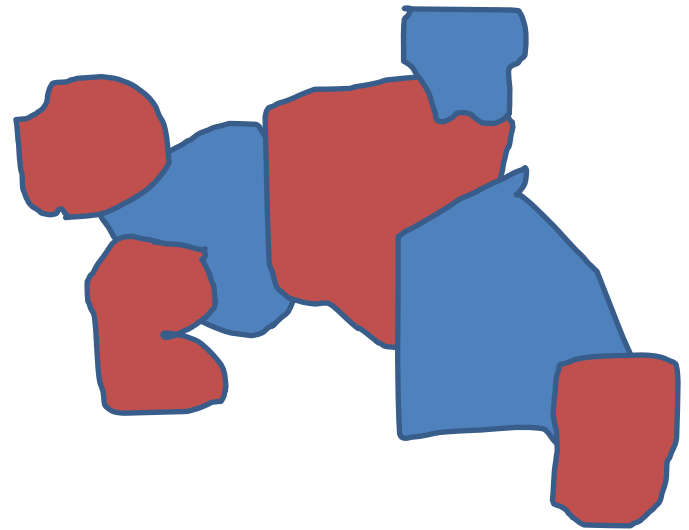
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# Colouring with 2 colours

- So checking if a map can be coloured with 2 colours is an easier problem than with three!
- And any map can be coloured with 4 colours.

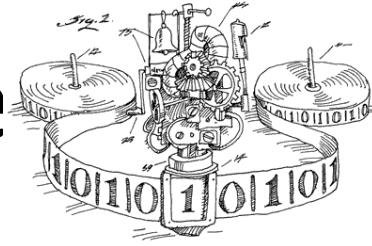


# Colouring maps

- If somebody gives you a coloured map, easy to check.
  - Check that there are 3 colours overall
  - Check that each country is different from its neighbours.
  - Done!
- Finding a colouring seems much harder...



# Polynomial-time computable



- Efficiently solvable:
  - On an input string of length  $n$
  - Produce a solution roughly in time at most
    - $n$ , or  $n^2$ , or  $n^3$ , or...  $n^{\text{const}}$ .
  - So a DNA string can be processed in about  $3.2 \times 10^9$  steps. At 1GHz, it is 3.2 seconds.
- Concept dates back to 1960s, Jack Edmonds, and also Alan Cobham.
  - Edmonds arguing why his “blossom algorithm” is better than what was known before.
- Checking that a given map is 3-colourable is polynomial time computable.
  - So is figuring out if a given map is 2-colourable.





# NP-completeness

- Is it possible to eliminate exhaustive search?
- **NP-completeness:** enough to answer for the problem of map colouring!
- A map is like a basket of apples.  
A map is colourable with 3 colours  
=  
There is a bad apple in the basket.
- Can “disguise” every problem that has efficiently checkable solutions as map colouring!
  - The concept of **NP-completeness** was invented by Stephen Cook (and independently Leonid Levin) in 1971
  - Made its way into popular culture, often as a synonym to “hard”... though we do not know for sure!



Stephen  
Cook



Leonid  
Levin

# Maps vs. teams

- Colouring a map with 3 colours is the same problem as splitting into three teams – in disguise.
- Call them TeamRed, TeamGreen and TeamYellow.
- Require that neighbours cannot be on the same team.
- If can split into three teams, colour countries by their team's colour.
- So if can split into teams efficiently, can colour maps just as efficiently.
- So splitting into two teams seems easier than into three!



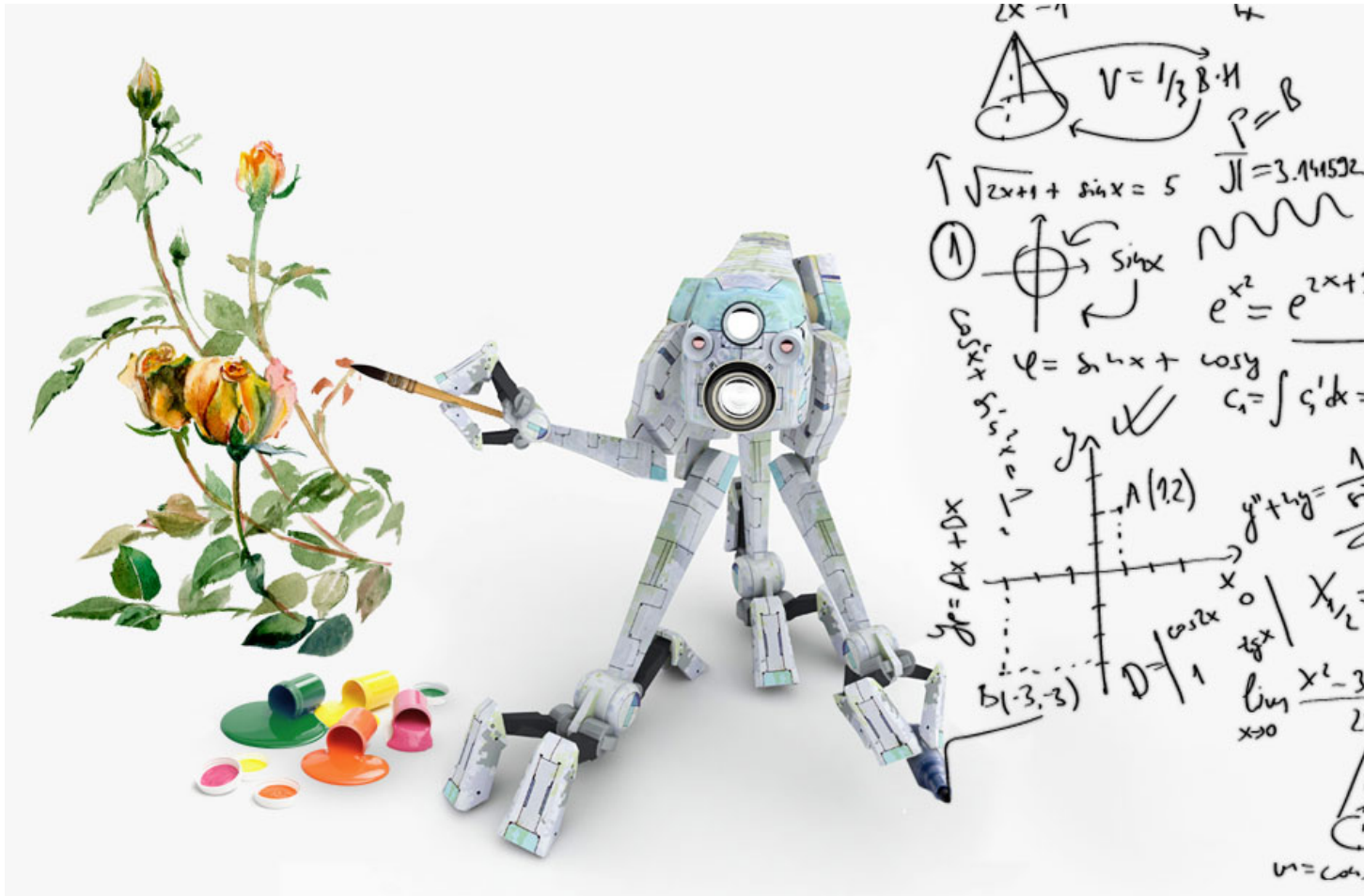


# P vs. NP

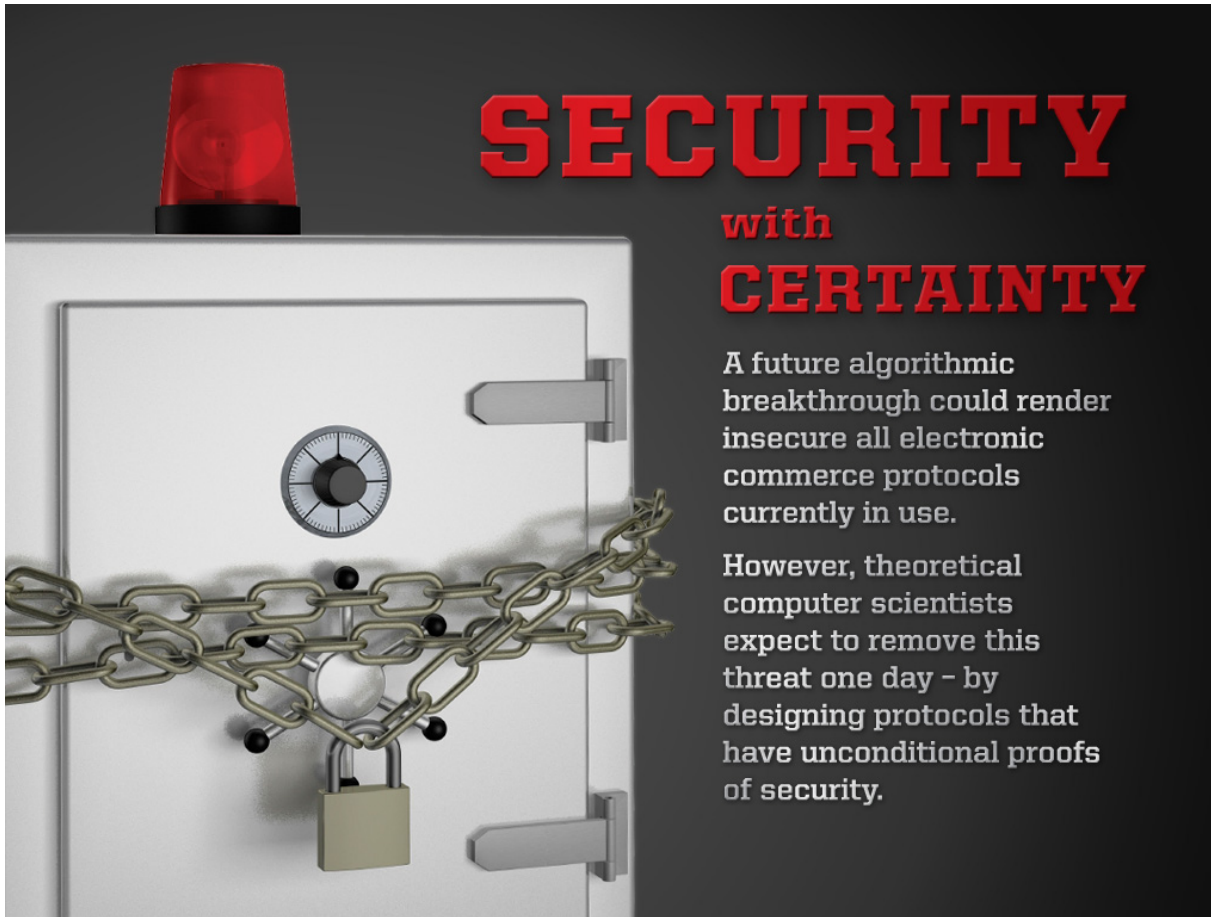
- If somebody finds a way to solve 3-colouring efficiently, then we will live in a very different world, where
  - Creativity and problem-solving are automated.
  - Not much security left on the internet.
  - Every theorem has a short proof...
- So most scientists believe that solving 3-coloring is impossible, but nobody so far can prove it.



# If P=NP... creativity is automated



# If $P=NP$ ... there is no security



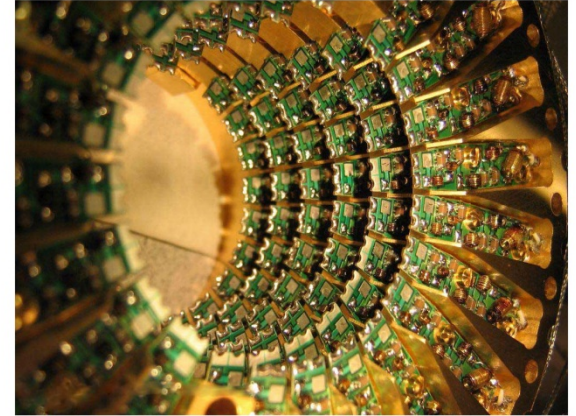
**SECURITY**  
with  
**CERTAINTY**

A future algorithmic breakthrough could render insecure all electronic commerce protocols currently in use.

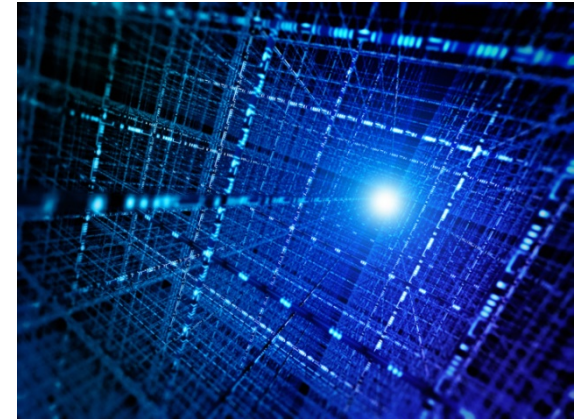
However, theoretical computer scientists expect to remove this threat one day - by designing protocols that have unconditional proofs of security.

# Quantum computers

- Can quantum computers colour maps efficiently?
- We don't know... but don't think so.
- Although they can factor numbers, which we do not know how to do on a usual computer fast.
- A real scalable quantum computer would require changing much of security on the internet.
  - RSA cryptosystem assumes factoring is hard.



Adiabatic Quantum Computer Component Array





# Exhaustive search beyond P vs. NP

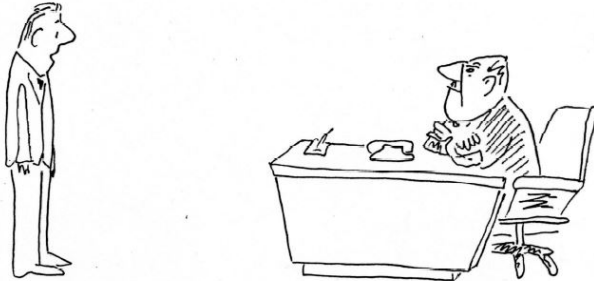


- Consider your friends on Facebook.
- Now look at the groups that they form
  - Belonging to the same Facebook group
  - Following the same person
  - Etc...
    - There can be many more groups than people.
- How hard is it to check if there are two groups that do not have any person in common?
  - Easy! Just compare any two groups.
    - Time to compare two groups times square of the number of groups.
  - Wait, but that's exhaustive search!... Can we do faster?..
    - If there are \*lots\* of groups, too slow in practice.
  - If so, can eliminate exhaustive search from map colouring...
    - Not quite solving P vs. NP, but getting there!





# P vs. NP (David Johnson's cartoons)



"I CAN'T SOLVE IT - I GUESS I'M JUST TOO DUMB."



"I CAN'T SOLVE IT - BECAUSE NO SOLUTION EXISTS!"



"I CAN'T SOLVE IT - BUT NEITHER CAN ALL THESE FAMOUS PEOPLE!"



WE MAY NOT BE ABLE TO SOLVE IT...  
BUT WE SURE CAN GET CLOSE!